

Access Control and Privacy Policies (8)

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Slides: KEATS (also homework is there)

Last Week

Andrew Secure RPC Protocol: A and B share a key private K_{AB} and want to identify each other

- A sends B : A, N_A
- B sends A : $\{N_A, K'_{AB}\}_{K_{AB}}$
- A sends B : $\{N_A\}_{K'_{AB}}$

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but most likely they are programs, which just follow some instructions
- indicates one “protocol run”, or session, which specifies some order in the communication
- there can be several sessions in parallel (think of wifi routers)

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A **reflection attack**: an intruder I impersonates B .

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Sounds stupid: "... answering a question with a counter question"

was originally developed at CMU for terminals to connect to workstations (e.g., file servers)

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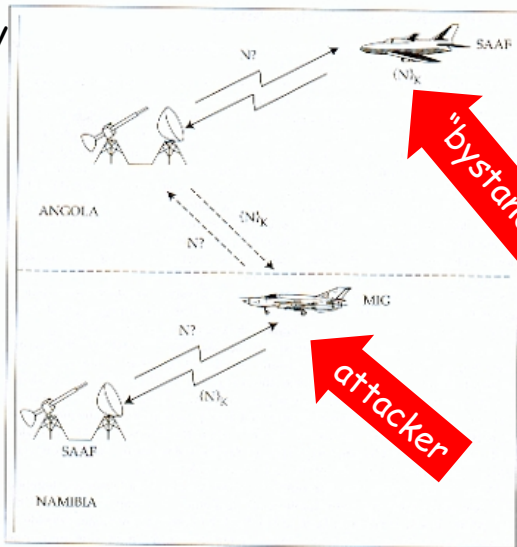


Figure 2.2 The MIG-in-the-middle attack.

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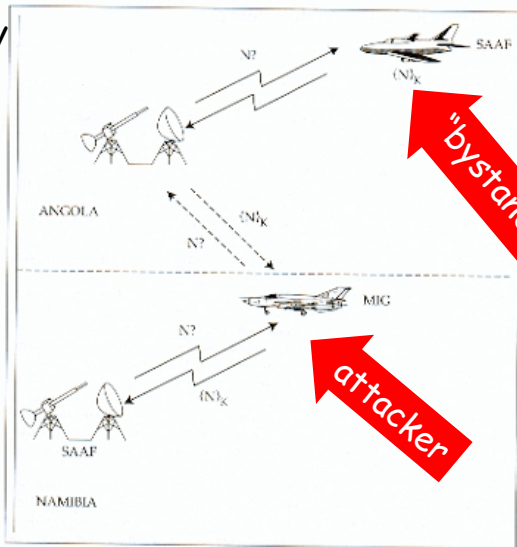


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IFF opened up a nice
side-channel attack

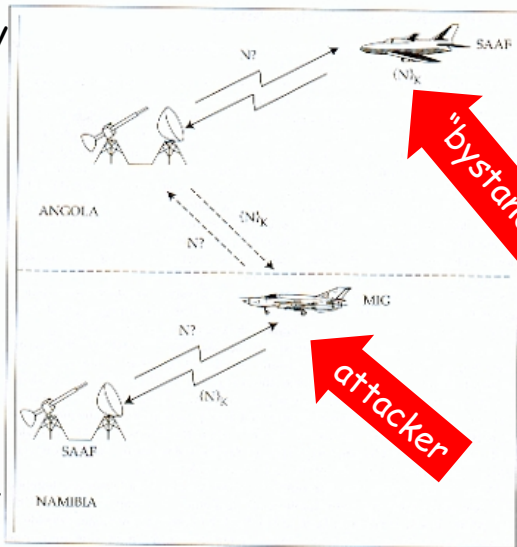


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Encryption to the Rescue?

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means you need to send separate "Hello" signals (bad), or worse share a single key between many entities

Protocol Attacks

- replay attacks
- reflection attacks
- man-in-the-middle attacks
- timing attacks
- parallel session attacks
- binding attacks (public key protocols)
- changing environment / changing assumptions

Replay Attacks

Schroeder-Needham protocol: exchange of a symmetric key with a trusted 3rd-party S :

$$A \rightarrow S : A, B, N_A$$

$$S \rightarrow A : \{N_A, B, K_{AB}, \{K_{AB}, A\}_{K_{BS}}\}_{K_{AS}}$$

$$A \rightarrow B : \{K_{AB}, A\}_{K_{BS}}$$

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at the end of the protocol both A and B should be in the possession of the secret key K_{AB} and know that the other principal has the key

Nonces

- 1 I generate a nonce (random number) and send it to you encrypted with a key we share
- 2 you increase it by one, encrypt it under a key I know and send it back to me

I can infer:

- you must have received my message
- you could only have generated your answer after I send you my initial message
- if only you and me know the key, the message must have come from you

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B believes it is following the correct protocol,
intruder I can form the correct response because
it knows K_{AB} and talks to B masquerading as A



"On the Internet, nobody knows you're a dog."

Replay Attacks

Andrew Secure RPC protocol: exchanging a new key between A and B

$$A \rightarrow B : A, \{N_A\}_{K_{AB}}$$

$$B \rightarrow A : \{N_A + 1, N_B\}_{K_{AB}}$$

$$A \rightarrow B : \{N_B + 1\}_{K_{AB}}$$

$$B \rightarrow A : \{K_{AB}^{new}, N_B^{new}\}_{K_{AB}}$$

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$$\begin{aligned}A &\rightarrow B : A, \{N_A\}_{K_{AB}} \\B &\rightarrow A : \{N_A + 1, N_B\}_{K_{AB}} \\A &\rightarrow B : \{N_B + 1\}_{K_{AB}} \\B &\rightarrow A : \{K_{AB}^{new}, N_B^{new}\}_{K_{AB}}\end{aligned}$$

Assume nonces are represented as bit-sequences of the same length

$$\begin{aligned}A &\rightarrow B : A, \{N_A\}_{K_{AB}} \\B &\rightarrow A : \{N_A + 1, N_B\}_{K_{AB}} \\A &\rightarrow I(B) : \{N_B + 1\}_{K_{AB}} \text{ intercepts} \\I(B) &\rightarrow A : \{N_A + 1, N_B\}_{K_{AB}} \text{ resend 2nd msg}\end{aligned}$$

Time-Stamps

The Schroeder-Needham protocol can be fixed by including a time-stamp (e.g., in Kerberos):

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but nothing is for free: then you need to synchronise time and possibly become victim to timing attacks

It can also be fixed by including another nonce:

$A \rightarrow B : A$

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Binding Attacks

with public-private keys it is important that the public key is **bound** to the right owner (verified by a certification authority CA)

$$A \rightarrow CA : A, B, N_A$$
$$CA \rightarrow A : CA, \{CA, A, N_A, K_B^{pub}\}_{K_A^{pub}}$$

A knows K_A^{priv} and can verify the message came from CA in response to A 's message and trusts K_B^{pub} is B 's public key

Binding Attacks

$A \rightarrow I(CA) : A, B, N_A$

$I(A) \rightarrow CA : A, I, N_A$

$CA \rightarrow I(A) : CA, \{CA, A, N_A, K_I^{pub}\}_{K_A^{pub}}$

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A now encrypts messages for B with the public key of I (which happily decrypts them with its private key)

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attacks because of changing environment

Changing Environment Attacks

- all protocols rely on some assumptions about the environment (e.g., cryptographic keys cannot be broken)

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- in the “good olden days” (1960/70) rail transport was cheap, so fraud was not worthwhile

Changing Environment Attacks

- all protocols rely on some assumptions about the environment (e.g., cryptographic keys cannot be broken)
- when it got expensive, some people bought cheaper monthly tickets for a suburban station and a nearby one, and one for the destination and a nearby one
- a large investment later all barriers were automatic and tickets could record state

Changing Environment Attacks

- all protocols rely on some assumptions about the environment (e.g., cryptographic keys cannot be broken)
- But suddenly the environment changed: rail transport got privatised creating many companies cheating each other
- revenue from monthly tickets was distributed according to a formula where the ticket was bought

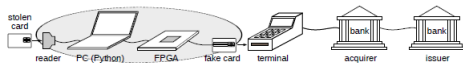
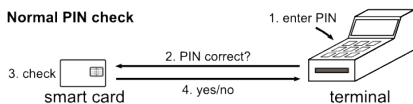
Changing Environment Attacks

- all protocols rely on some assumptions about the environment (e.g., cryptographic keys cannot be broken)
- apart from bad outsiders (passengers) you also had bad insiders (rail companies)
- chaos and litigation ensued

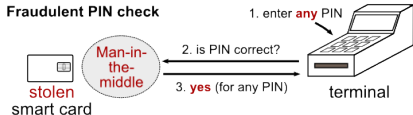
A Man-in-the-middle attack in real life:

- the card only says yes or no to the terminal if the PIN is correct
- trick the card in thinking transaction is verified by signature
- trick the terminal in thinking the transaction was verified by PIN

Normal PIN check



Fraudulent PIN check



Problems with EMV

- it is a wrapper for many protocols
- specification by consensus (resulted unmanageable complexity)
- its specification is 700 pages in English plus 2000+ pages for testing, additionally some further parts are secret
- other attacks have been found
- one solution might be to require always online verification of the PIN with the bank

Good Practices

- explicit principles (you authenticate all data you might rely on)
- the one who can fix a system should also be liable for the losses

Privacy et al

Some terminology:

- **secrecy** is the mechanism used to limit the number of principals with access to information (eg, cryptography or access controls)
- **confidentiality** is the obligation to protect the secrets of other people or organizations (secrecy for the benefit of an organisation)
- **anonymity** is the ability to leave no evidence of an activity (eg, sharing a secret)
- **privacy** is the ability or right to protect your personal secrets (secrecy for the benefit of an individual)

Privacy vs Anonymity

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But privacy?

"You have zero privacy anyway. Get over it."

Scott Mcnealy (CEO of Sun)

If you have nothing to hide, you have nothing to fear.

Privacy

private data can be often used against me

- if my location data becomes public, thieves will switch off their phones and help themselves in my home
- if supermarkets can build a profile of what I buy, they can use it to their advantage (banks - mortgages)
- my employer might not like my opinions

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