Access Control and Privacy Policies (9)

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Slides: KEATS (also homework is there)

Review: Proofs

Modify the formula

$$P \; \textit{controls} \; \textit{Permitted}(O, \textit{write})$$

using security levels so that it satisfies the *write rule* from the *Bell-LaPadula* access policy.

Do the same again, but satisfy the *write rule* from the *Biba* access policy.

Review: Proofs

Assume two security levels S and TS, which are ordered so that $\boldsymbol{slev}(S) < \boldsymbol{slev}(TS)$. Assume further the substitution rules

$$\frac{\Gamma \vdash \boldsymbol{slev}(\boldsymbol{P}) = \boldsymbol{l}_1 \quad \Gamma \vdash \boldsymbol{slev}(\boldsymbol{Q}) = \boldsymbol{l}_2 \quad \Gamma \vdash \boldsymbol{l}_1 < \boldsymbol{l}_2}{\Gamma \vdash \boldsymbol{slev}(\boldsymbol{P}) < \boldsymbol{slev}(\boldsymbol{Q})}$$

$$\frac{\Gamma \vdash \boldsymbol{slev}(\boldsymbol{P}) = \boldsymbol{l} \quad \Gamma \vdash \boldsymbol{slev}(\boldsymbol{Q}) = \boldsymbol{l}}{\Gamma \vdash \boldsymbol{slev}(\boldsymbol{P}) = \boldsymbol{slev}(\boldsymbol{Q})}$$

Review: Proofs

Let Γ be the set containing the following six formulas

$$slev(S) < slev(TS)$$
$$slev(Agent) = slev(TS)$$
$$slev(File_1) = slev(S)$$
$$slev(File_2) = slev(TS)$$
$$\forall O.\ slev(O) < slev(Agent) \Rightarrow (Agent\ controls\ Permitted(O, read))$$
$$\forall O.\ slev(O) = slev(Agent) \Rightarrow (Agent\ controls\ Permitted(O, read))$$

Using the inference rules of access-control logic and the substitution rules shown above, give proofs for the two judgements

$$\Gamma \vdash (Agent\ says\ Permitted(File_1, read)) \Rightarrow Permitted(File_1, read)$$

Checking Solutions

How can you check somebody's solution without revealing the solution?

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- folio

  "an individual leaf of paper or parchment, either loose as one of a series or forming part of a bound volume, which is numbered on the recto or front side only."

Checking Solutions

How can you check somebody's solution without revealing the solution?

Alice and Bob solve crosswords. Alice knows the answer for 21D (folio) but doesn't want to tell Bob.

You use an English dictionary:

- folio $\xrightarrow{1}$ individual
  "a single human being as distinct from a group"

Checking Solutions

How can you check somebody's solution without revealing the solution?

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You use an English dictionary:

- folio $\xrightarrow{1}$ individual $\xrightarrow{2}$ human

  "relating to or characteristic of humankind"

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this is essentially a hash function...but Bob can only check once he has also found the solution

Zero-Knowledge Proofs

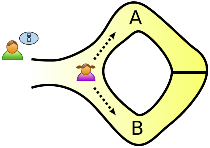
Two remarkable properties of Zero-Knowledge Proofs:

- Alice only reveals the fact that she knows a secret, not the secret itself (meaning she can convince Bob that she knows the secret).
- Having been convinced, Bob cannot use the evidence in order to convince Carol that Alice knows the secret.

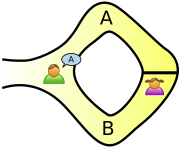
The Idea

The Alibaba cave:

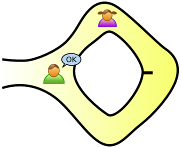
1.



2.



3.



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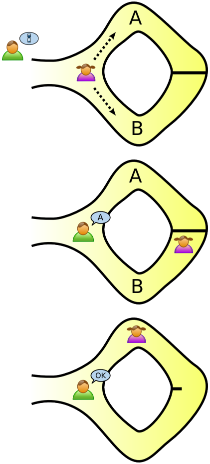

Even if Bob has a hidden camera, a recording will not be convincing to anyone else (Alice and Bob could have made it all up).

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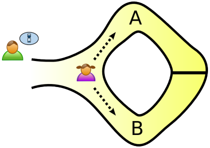

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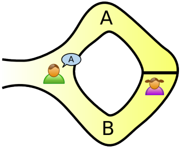

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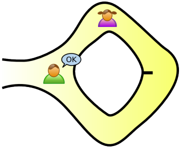


Even worse, an observer present at the experiment would not be convinced.

Applications of ZKPs

- authentication, where one party wants to prove its identity to a second party via some secret information, but doesn't want the second party to learn anything about this secret

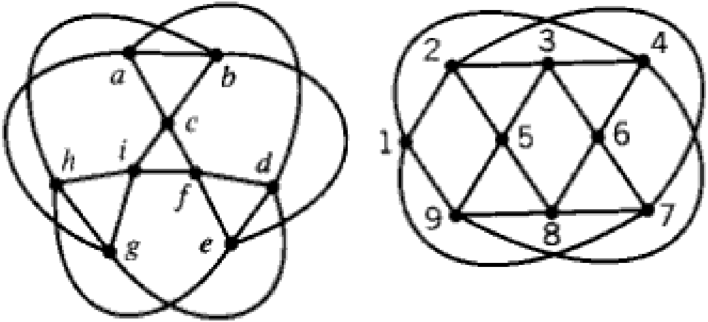
- to enforce honest behaviour while maintaining privacy: the idea is to force a user to prove, using a zero-knowledge proof, that its behaviour is correct according to the protocol

Central Properties

Zero-knowledge proof protocols should satisfy:

- **Completeness** If Alice knows the secret, Bob accepts Alice "proof" for sure.
- **Soundness** If Alice does not know the secret, Bob accepts her "proof" with a very small probability.

Graph Isomorphism



Finding an isomorphism between two graphs is an NP complete problem.

Graph Isomorphism Protocol

Alice starts with knowing an isomorphism σ between graphs G_1 and G_2

1. Alice generates an isomorphic graph H which she sends to Bob
2. Bob asks either for an isomorphism between G_1 and H, or G_2 and H
3. Alice and Bob repeat this procedure n times

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these are called commitment algorithms

Graph Isomorphism Protocol (2)

If Alice knows the isomorphism, she can always calculate σ.

If she doesn't, she can only correctly respond if Bob's choice of index is the same as the one she used to form H. The probability of this happening is $\frac{1}{2}$, so after n rounds the probability of her always responding correctly is only $\frac{1}{2}^n$.

Graph Isomorphism Protocol (3)

Why is the GI-protocol zero-knowledge?

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Why is the GI-protocol zero-knowledge?

A: We can generate a fake transcript of a conversation, which cannot be distinguished from a "real" conversation.

Anything Bob can compute using the information obtained from the transcript can be computed using only a forged transcript and therefore participation in such a communication does not increase Bob's capability to perform any computation.

Non-Interactive ZKPs

This is amazing: Alison can publish some data that contains no data about her secret, but this data can be used to convince anyone of the secret's existence.

Non-Interactive ZKPs (2)

Alice starts with knowing an isomorphism σ between graphs G_1 and G_2

1. Alice generates n isomorphic graphs $H_{1..n}$ which she makes public
2. she feeds the $H_{1..n}$ into a hashing function (she has no control over what the output will be)
3. Alice takes the first n bits of the output: whenever output is 0, she shows an isomorphism with G_1 ; for 1 she shows an isomorphism with G_2

Problems of ZKPs

- “grand chess master problem”
(person in the middle)

- Alice can have multiple identities; once she committed a fraud she stops using one

Other Methods for ZKPs

- modular logarithms (you can

Random Number Generators