## A Formalisation of an Access Control Framework





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#### **Access Control**

 perhaps most known are Unix-style access control systems (root super-user, setuid mechanism)

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more fine-grained access control is provided by

- SELinux (security enhanced Linux devloped by the NSA; mandatory access control system)
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## Operations in the OS

using Paulson's inductive method a **state of the system** is a **trace**, a list of events (system calls):

$$[\boldsymbol{e}_1,\ldots,\boldsymbol{e}_2]$$

we need to restrict the traces to **valid traces**:

```
valid [] valid s admissible s e granted s e valid (e::s)
```

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 $p \in current\_procs \ s$   $p' \notin current\_procs \ s$   $admissible \ s \ (Clone \ p \ p')$ 

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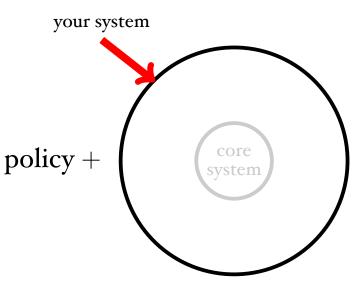


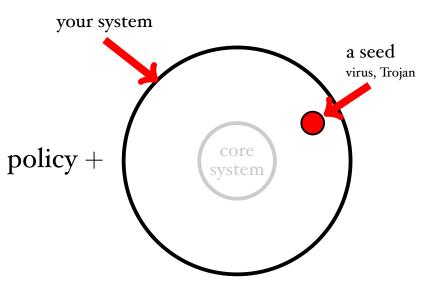
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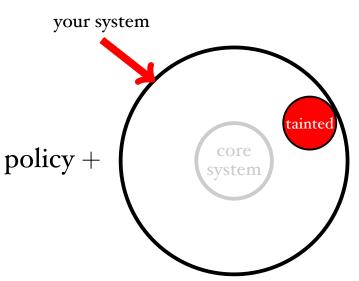
is\_current\_role  $s \ p \ r$  is\_file\_type  $s \ f \ t$   $(r, t, Execute) \in permissions$   $granted \ s \ (Execute \ p \ f)$ 

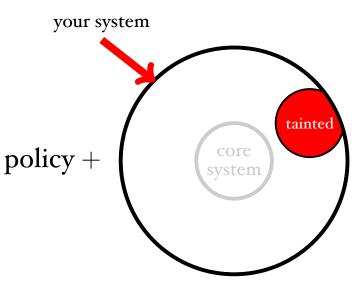
## **Design of AC-Policies**

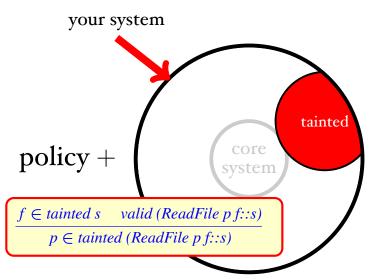
"what you specify is what you get but not necessarily what you want..."

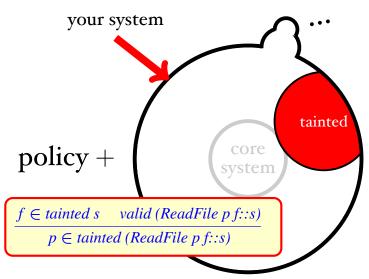












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- working purely on *static* policies also does not work – because of over approximation
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- our solution: take a middle ground and record precisely the information of the initial state, but be less precise about every newly created object.

#### Results about our Test

• we can show that the objects (files, processes, ...) we need to consider are only finite (at some point it does not matter if we create another *bin*-file)

#### Thm (Soundness)

If our test says an object is taintable, then it is taintable in the OS, and we produce a sequence of events that show how it can be tainted.

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#### Thm (Completeness)

If an object is taintable in the OS and *undeletable*\*, then our test will find out that it is taintable.

\* an object is *undeleteable* if it exists in the initial state, but there exists no valid state in which it could have been deleted

#### Why the Restriction?

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- after that a proces with the same ID gets re-created by cloning an untainted process
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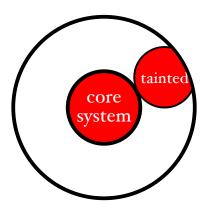
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Is this a serious restriction? We think not ...

## **Core System**

Admins usually ask whether their policy is strong enough to protect their core system?



core system files are typically undeletable

#### **Conclusion**

- we considered the Role-Compatibility Model used for securing the Apache Server
   13 events, 13 rules for OS admissibility, 14 rules for RC-granting, 10 rules for tainted
- we can scale this to SELinux
  more fine-grainded OS events (inodes, hard-links, shared memory, ...)
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  - 22 events, 22 admisibility, 22 granting, 15 taintable
- hard sell to Ott (who designed the RC-model)
- hard sell to the community working on access control (beyond good science)

# Thanks! Questions?