

A lemma which might be true, but can also be false, is as follows:

- If
- (1) $v_1 \succ_{der\ c\ r} v_2$,
 - (2) $\vdash v_1 : der\ c\ r$, and
 - (3) $\vdash v_2 : der\ c\ r$ holds,
- then $inj\ r\ c\ v_1 \succ_r inj\ r\ c\ v_2$ also holds.

It essentially states that if one value v_1 is bigger than v_2 then this ordering is preserved under injections. This is proved by induction (on the definition of der ... this is very similar to an induction on r).

The case that is still unproved is the sequence case where we assume $r = r_1 \cdot r_2$ and also r_1 being nullable. The derivative $der\ c\ r$ is then

$$der\ c\ r = ((der\ c\ r_1) \cdot r_2) + (der\ c\ r_2)$$

or without the parentheses

$$der\ c\ r = (der\ c\ r_1) \cdot r_2 + der\ c\ r_2$$

In this case the assumptions are

- (a) $v_1 \succ_{(der\ c\ r_1) \cdot r_2 + der\ c\ r_2} v_2$
- (b) $\vdash v_1 : (der\ c\ r_1) \cdot r_2 + der\ c\ r_2$
- (c) $\vdash v_2 : (der\ c\ r_1) \cdot r_2 + der\ c\ r_2$
- (d) $nullable(r_1)$

The induction hypotheses are

- (IH1) $\forall v_1 v_2. v_1 \succ_{der\ c\ r_1} v_2 \wedge \vdash v_1 : der\ c\ r_1 \wedge \vdash v_2 : der\ c\ r_1$
 $\longrightarrow inj\ r_1\ c\ v_1 \succ_{r_1} inj\ r_1\ c\ v_2$
- (IH2) $\forall v_1 v_2. v_1 \succ_{der\ c\ r_2} v_2 \wedge \vdash v_2 : der\ c\ r_2 \wedge \vdash v_2 : der\ c\ r_2$
 $\longrightarrow inj\ r_2\ c\ v_1 \succ_{r_2} inj\ r_2\ c\ v_2$

The goal is

$$(goal) \quad inj\ (r_1 \cdot r_2)\ c\ v_1 \succ_{r_1 \cdot r_2} inj\ (r_1 \cdot r_2)\ c\ v_2$$

If we analyse how (a) could have arisen (that is make a case distinction), then we will find four cases:

- LL $v_1 = Left(w_1), v_2 = Left(w_2)$
- LR $v_1 = Left(w_1), v_2 = Right(w_2)$
- RL $v_1 = Right(w_1), v_2 = Left(w_2)$
- RR $v_1 = Right(w_1), v_2 = Right(w_2)$

We have to establish our goal in all four cases.

Case LR

The corresponding rule (instantiated) is:

$$\frac{\text{len } |w_1| \geq \text{len } |w_2|}{\text{Left}(w_1) \succ_{(\text{der } c \ r_1) \cdot r_2 + \text{der } c \ r_2} \text{Right}(w_2)}$$

This means we can also assume in this case

$$(e) \quad \text{len } |w_1| \geq \text{len } |w_2|$$

which is the premise of the rule above. Instantiating v_1 and v_2 in the assumptions (b) and (c) gives us

$$\begin{aligned} (b^*) \quad & \vdash \text{Left}(w_1) : (\text{der } c \ r_1) \cdot r_2 + \text{der } c \ r_2 \\ (c^*) \quad & \vdash \text{Right}(w_2) : (\text{der } c \ r_1) \cdot r_2 + \text{der } c \ r_2 \end{aligned}$$

Since these are assumptions, we can further analyse how they could have arisen according to the rules of $\vdash _ : _$. This gives us two new assumptions

$$\begin{aligned} (b^{**}) \quad & \vdash w_1 : (\text{der } c \ r_1) \cdot r_2 \\ (c^{**}) \quad & \vdash w_2 : \text{der } c \ r_2 \end{aligned}$$

Looking at (b^{**}) we can further analyse how this judgement could have arisen. This tells us that w_1 must have been a sequence, say $u_1 \cdot u_2$, with

$$\begin{aligned} (b^{***}) \quad & \vdash u_1 : \text{der } c \ r_1 \\ & \vdash u_2 : r_2 \end{aligned}$$

Instantiating the goal means we need to prove

$$\text{inj } (r_1 \cdot r_2) \ c \ (\text{Left}(u_1 \cdot u_2)) \succ_{r_1 \cdot r_2} \text{inj } (r_1 \cdot r_2) \ c \ (\text{Right}(w_2))$$

We can simplify this according to the rules of *inj*:

$$(\text{inj } r_1 \ c \ u_1) \cdot u_2 \succ_{r_1 \cdot r_2} (\text{mkeps } r_1) \cdot (\text{inj } r_2 \ c \ w_2)$$

This is what we need to prove.

Case RL

The corresponding rule (instantiated) is:

$$\frac{\text{len } |w_1| > \text{len } |w_2|}{\text{Right}(w_1) \succ_{(\text{der } c \ r_1) \cdot r_2 + \text{der } c \ r_2} \text{Left}(w_2)}$$