Compilers and Formal Languages (1)

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The Goal of this Course

Write A Compiler



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lexer input: a string

"read(n);"

lexer output: a sequence of tokens
 key(read); lpar; id(n); rpar; semi









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The subject is quite old

- Turing Machines, 1936
- Regular Expressions, 1956
- The first compiler for COBOL, 1957 (Grace Hopper)
- But surprisingly research papers are still published nowadays



Grace Hopper

(she made it to David Letterman's Tonight Show,

http://www.youtube.com/watch?v=aZOxtURhfEU)





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Lectures 1 - 5

transforming strings into structured data

Lexing based on regular expressions (recognising "words")

Parsing

(recognising "sentences")



Stone of Rosetta

Familiar Regular Expr.

- re* matches o or more times matches 1 or more times re+ re? matches o or I times re{n} matches exactly n number of times re{n,m} matches at least n and at most m times matches any single character inside the brackets [...] [^...] matches any single character not inside the brackets character ranges a-zA-Z ١d matches digits; equivalent to [0-9] matches every character except newline
- (re) groups regular expressions and remembers the matched text

Today

• While the ultimate goal is to implement a small compiler (a really small one for the JVM)...

Let's start with:

- a web-crawler
- an email harvester
- (a web-scraper)



- given an URL, read the corresponding webpage
- extract all links from it
- I call the web-crawler again for all these links



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- if not possible print, out a problem
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(we need a bound for the number of recursive calls) (the purpose is to check all links on my own webpage)





A simple Scala function for reading webpages:

```
import io.Source
```

```
def get_page(url: String) : String = {
   Source.fromURL(url).take(10000).mkString
}
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get_page(""""http://www.inf.kcl.ac.uk/staff/urbanc/""")

A slightly more complicated version for handling errors:

```
def get_page(url: String) : String = {
   Try(Source.fromURL(url).take(10000).mkString).
     getOrElse { println(s" Problem with: $url"); ""}
}
```

A Regular Expression

• ... is a pattern or template for specifying strings

"https?://[^"]*"

matches for example "http://www.foobar.com" "https://www.tls.org"

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"""""https?://[^"]*"""".r

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Finding Operations

rexp.findAllIn(string)

returns a list of all (sub)strings that match the regular expression

rexp.findFirstIn(string) returns either

- None if no (sub)string matches or
- Some(s) with the first (sub)string

val http_pattern = """"https?://[^"]*"""".r

def unquote(s: String) = s.drop(1).dropRight(1)

```
def get_all_URLs(page: String) : Set[String] =
    http_pattern.findAllIn(page).map(unquote).toSet
```

```
def crawl(url: String, n: Int) : Unit = {
    if (n == 0) ()
    else {
        println(s"Visiting: $n $url")
        for (u <- get_all_URLs(get_page(url))) crawl(u, n - 1)
    }
}</pre>
```

```
crawl(some_start_URL, 2)
```

```
A version that only crawls links in "my" domain:
val my urls = """urbanc""".r
def crawl(url: String, n: Int) : Unit = {
  if(n == 0)()
  else if (my urls.findFirstIn(url) == None) {
    println(s"Visiting: $n $url")
    get page(url); ()
  else {
    println(s"Visiting: $n $url")
    for (u <- get_all_URLs(get_page(url))) crawl(u, n - 1)</pre>
}
```

A little email harvester:

```
val http pattern = """"https?://[^"]*"""".r
val email pattern =
  """([a-z0-9_\.-]+)@([\da-z\.-]+)\.([a-z\.]{2,6})""".r
def print str(s: String) =
  if (s == "") () else println(s)
def crawl(url: String, n: Int) : Unit = {
  if(n == 0)()
  else {
    println(s"Visiting: $n $url")
    val page = get page(url)
    print str(email_pattern.findAllIn(page).mkString("\n"))
    for (u <- get all URLs(page).par) crawl(u, n - 1)</pre>
  }
```

http://net.tutsplus.com/tutorials/other/8-regular-expressions-you-should-know/

Regular Expressions

Their inductive definition:



```
Th

abstract class Rexp

case object ZERO extends Rexp

case object ONE extends Rexp

case class CHAR(c: Char) extends Rexp

case class ALT(r1: Rexp, r2: Rexp) extends Rexp

case class SEQ(r1: Rexp, r2: Rexp) extends Rexp

case class STAR(r: Rexp) extends Rexp
```

r :::	= 0	null
	I	empty string / "" / []
	С	character
	$r_{\rm I}+r_2$	alternative / choice
	$r_{I} \cdot r_{2}$	sequence
	<i>r</i> *	star (zero or more)



... are lists of characters. For example "hello"

[b, e, l, l, o] or just *hello*

the empty string: [] or ""

the concatenation of two strings:

 $s_1 @ s_2$

foo @ bar = foobar, baz @ [] = baz

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Languages, Strings

- **Strings** are lists of characters, for example [], *abc* (Pattern match: *c*::*s*)
- A **language** is a set of strings, for example {[], *bello*, *foobar*, *a*, *abc*}
- **Concatenation** of strings and languages

foo @ bar = foobar $A @ B \stackrel{\text{def}}{=} \{s_1 @ s_2 \mid s_1 \in A \land s_2 \in B\}$

The Meaning of a **Regular Expression** $L(\mathbf{0}) \stackrel{\text{def}}{=} \{\}$ $L(\mathbf{I}) \stackrel{\text{def}}{=} \{[]\}$ $L(c) \stackrel{\text{def}}{=} \{[c]\}$ $L(r_{I}+r_{2}) \stackrel{\text{def}}{=} L(r_{I}) \cup L(r_{2})$ $L(\mathbf{r}_{\mathrm{I}} \cdot \mathbf{r}_{2}) \stackrel{\text{def}}{=} \{ s_{\mathrm{I}} @ s_{2} \mid s_{\mathrm{I}} \in L(\mathbf{r}_{\mathrm{I}}) \land s_{2} \in L(\mathbf{r}_{2}) \}$ $L(\mathbf{r}^*) \stackrel{\text{def}}{=}$

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 $\begin{array}{rcl} L(r)^{\circ} & \stackrel{\text{def}}{=} & \{[]\} \\ L(r)^{n+1} & \stackrel{\text{def}}{=} & L(r) @ L(r)^{n} & \text{(append on sets)} \\ & & \{s_{\mathrm{I}} @ s_{2} \mid s_{\mathrm{I}} \in L(r) \land s_{2} \in L(r)^{n}\} \end{array}$

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The Meaning of Matching

```
A regular expression r matches a string s provided
s \in L(r)
```

...and the point of the next lecture is to decide this problem as fast as possible (unlike Python, Ruby, Java)

Written Exam

- Accounts for 80%.
- You will understand the question "*Is this relevant for the exam?*" is very demotivating for the lecturer!
- Deal: Whatever is in the homework (and is not marked "*optional*") is relevant for the exam.
- Each lecture has also a handout. There are also handouts about notation and Scala.

Coursework

• Accounts for 20%. Two strands. Choose one!

Strand 1

- four programming tasks:
 - matcher (4%, 19.10.)
 - lexer (5%, 03.11.)
 - parser (5%, 23.11.)
 - compiler (6%, 7.12.)

Strand 2

- one task: prove the correctness of a regular expression matcher in the Isabelle theorem prover
- 20%, submission 7.12.
- Solving more than one strand will **not** give you more marks.

Lecture Capture

• Hope it works...

Lecture Capture

- Hope it works...
- It is important to use lecture capture wisely:
 - Lecture recordings are a study and revision aid.
 - Statistically, there is a clear and direct link between attendance and attainment: Students who do not attend lectures, do less well in exams.
- Attending a lecture is more than watching it online if you do not attend, you miss out!



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