

# Compilers and Formal Languages

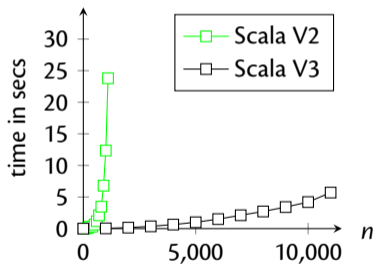
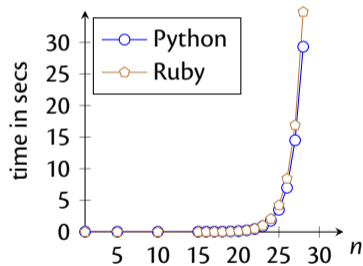
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Slides & Progs: KEATS (also homework is there)

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# Let's Implement an Efficient Regular Expression Matcher

Graphs:  $a^{?{n}} \cdot a^{n}$  and strings  $\underbrace{a \dots a}_n$



In the handouts is a similar graph for  $(a^*)^* \cdot b$  and Java 8, JavaScript and Python.

# (Basic) Regular Expressions

Their inductive definition:

$r ::=$	<b>0</b>	nothing
	<b>1</b>	empty string / "" / []
	$c$	character
	$r_1 + r_2$	alternative / choice
	$r_1 \cdot r_2$	sequence
	$r^*$	star (zero or more)

# When Are Two Regular Expressions Equivalent?

Two regular expressions  $r_1$  and  $r_2$  are **equivalent** provided:

$$r_1 \equiv r_2 \stackrel{\text{def}}{=} L(r_1) = L(r_2)$$

# Some Concrete Equivalences

$$(a + b) + c \equiv a + (b + c)$$

$$a + a \equiv a$$

$$a + b \equiv b + a$$

$$(a \cdot b) \cdot c \equiv a \cdot (b \cdot c)$$

$$c \cdot (a + b) \equiv (c \cdot a) + (c \cdot b)$$

# Some Concrete Equivalences

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$$(a \cdot b) \cdot c \equiv a \cdot (b \cdot c)$$

$$c \cdot (a + b) \equiv (c \cdot a) + (c \cdot b)$$

$$a \cdot a \not\equiv a$$

$$a + (b \cdot c) \not\equiv (a + b) \cdot (a + c)$$

# Some Corner Cases

$$\begin{array}{l} a \cdot 0 \not\equiv a \\ a + 1 \not\equiv a \\ 1 \equiv 0^* \\ 1^* \equiv 1 \\ 0^* \not\equiv 0 \end{array}$$

# Some Simplification Rules

$$r + 0 \equiv r$$

$$0 + r \equiv r$$

$$r \cdot 1 \equiv r$$

$$1 \cdot r \equiv r$$

$$r \cdot 0 \equiv 0$$

$$0 \cdot r \equiv 0$$

$$r + r \equiv r$$



# Simplification Example

$$\begin{aligned}((\mathbf{1} \cdot b) + \mathbf{0}) \cdot r &\Rightarrow ((\underline{\mathbf{1} \cdot b}) + \mathbf{0}) \cdot r \\ &= (\underline{b + \mathbf{0}}) \cdot r \\ &= b \cdot r\end{aligned}$$

# Simplification Example

$$\begin{aligned}((\mathbf{0} \cdot b) + \mathbf{0}) \cdot r &\Rightarrow ((\underline{\mathbf{0} \cdot b}) + \mathbf{0}) \cdot r \\ &= (\underline{\mathbf{0} + \mathbf{0}}) \cdot r \\ &= \mathbf{0} \cdot r \\ &= \mathbf{0}\end{aligned}$$

# Semantic Derivative

- The **Semantic Derivative** of a language w.r.t. to a character  $c$ :

$$\text{Der } c A \stackrel{\text{def}}{=} \{s \mid c :: s \in A\}$$

For  $A = \{\text{foo}, \text{bar}, \text{frak}\}$  then

$$\text{Der } f A = \{\text{oo}, \text{rak}\}$$

$$\text{Der } b A = \{\text{ar}\}$$

$$\text{Der } a A = \{\}$$

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$$\text{Der } b A = \{\text{ar}\}$$

$$\text{Der } a A = \{\}$$

We can extend this definition to strings

$$\text{Der } s A = \{s' \mid s @ s' \in A\}$$

# The Specification for Matching

A regular expression  $r$  matches a string  $s$  provided:

$$s \in L(r)$$

...and the point of the this lecture is to decide this problem as fast as possible (unlike Python, Ruby, Java etc)

# Brzowski's Algorithm (1)

...whether a regular expression can match the empty string:

$$\text{nullable}(\mathbf{0}) \stackrel{\text{def}}{=} \text{false}$$

$$\text{nullable}(\mathbf{1}) \stackrel{\text{def}}{=} \text{true}$$

$$\text{nullable}(c) \stackrel{\text{def}}{=} \text{false}$$

$$\text{nullable}(r_1 + r_2) \stackrel{\text{def}}{=} \text{nullable}(r_1) \vee \text{nullable}(r_2)$$

$$\text{nullable}(r_1 \cdot r_2) \stackrel{\text{def}}{=} \text{nullable}(r_1) \wedge \text{nullable}(r_2)$$

$$\text{nullable}(r^*) \stackrel{\text{def}}{=} \text{true}$$

# The Derivative of a Rexp

If  $r$  matches the string  $c::s$ , what is a regular expression that matches just  $s$ ?

$der\ c\ r$  gives the answer, Brzozowski 1964

# The Derivative of a Rexp

$$\text{der } c(0) \stackrel{\text{def}}{=} 0$$

$$\text{der } c(1) \stackrel{\text{def}}{=} 0$$

$$\text{der } c(d) \stackrel{\text{def}}{=} \text{if } c = d \text{ then } 1 \text{ else } 0$$

$$\text{der } c(r_1 + r_2) \stackrel{\text{def}}{=} \text{der } c r_1 + \text{der } c r_2$$

$$\text{der } c(r_1 \cdot r_2) \stackrel{\text{def}}{=} \text{if nullable}(r_1) \\ \text{then } (\text{der } c r_1) \cdot r_2 + \text{der } c r_2 \\ \text{else } (\text{der } c r_1) \cdot r_2$$

$$\text{der } c(r^*) \stackrel{\text{def}}{=} (\text{der } c r) \cdot (r^*)$$



# The Derivative of a Rexp

$$\text{der } c \text{ (0)} \stackrel{\text{def}}{=} \mathbf{0}$$

$$\text{der } c \text{ (1)} \stackrel{\text{def}}{=} \mathbf{0}$$

$$\text{der } c \text{ (} d \text{)} \stackrel{\text{def}}{=} \text{if } c = d \text{ then } \mathbf{1} \text{ else } \mathbf{0}$$

$$\text{der } c \text{ (} r_1 + r_2 \text{)} \stackrel{\text{def}}{=} \text{der } c \text{ } r_1 + \text{der } c \text{ } r_2$$

$$\text{der } c \text{ (} r_1 \cdot r_2 \text{)} \stackrel{\text{def}}{=} \text{if nullable}(r_1) \\ \text{then } (\text{der } c \text{ } r_1) \cdot r_2 + \text{der } c \text{ } r_2 \\ \text{else } (\text{der } c \text{ } r_1) \cdot r_2$$

$$\text{der } c \text{ (} r^* \text{)} \stackrel{\text{def}}{=} (\text{der } c \text{ } r) \cdot (r^*)$$

$$\text{ders } [] \text{ } r \stackrel{\text{def}}{=} r$$

$$\text{ders } (c :: s) \text{ } r \stackrel{\text{def}}{=} \text{ders } s \text{ (der } c \text{ } r)$$

# Examples

Given  $r \stackrel{\text{def}}{=} ((a \cdot b) + b)^*$  what is

$\text{der } a \ r = ?$

$\text{der } b \ r = ?$

$\text{der } c \ r = ?$

# Derivative Example

Given  $r \stackrel{\text{def}}{=} ((a \cdot b) + b)^*$  what is

$$\begin{aligned} \text{der } a ((a \cdot b) + b)^* &\Rightarrow \text{der } a \underline{((a \cdot b) + b)^*} \\ &= (\text{der } a \underline{((a \cdot b) + b)}) \cdot r \\ &= ((\text{der } a \underline{a \cdot b}) + (\text{der } a b)) \cdot r \\ &= (((\text{der } a \underline{a}) \cdot b) + (\text{der } a b)) \cdot r \\ &= ((\mathbf{1} \cdot b) + (\text{der } a \underline{b})) \cdot r \\ &= ((\mathbf{1} \cdot b) + \mathbf{0}) \cdot r \end{aligned}$$

# The Brzozowski Algorithm

$$\text{matcher } r s \stackrel{\text{def}}{=} \text{nullable}(\text{ders } s r)$$

# Brzowski: An Example

Does  $r_1$  match  $abc$ ?

Step 1: build derivative of  $a$  and  $r_1$  ( $r_2 = \text{der } a \ r_1$ )

Step 2: build derivative of  $b$  and  $r_2$  ( $r_3 = \text{der } b \ r_2$ )

Step 3: build derivative of  $c$  and  $r_3$  ( $r_4 = \text{der } c \ r_3$ )

Step 4: the string is exhausted: ( $\text{nullable}(r_4)$ )

test whether  $r_4$  can recognise  
the empty string

Output: result of the test  
 $\Rightarrow$   $\text{true}$  or  $\text{false}$

# The Idea of the Algorithm

If we want to recognise the string  $abc$  with regular expression  $r_1$  then

①  $Der a(L(r_1))$

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- 2  $Der b (Der a (L(r_1)))$

# The Idea of the Algorithm

If we want to recognise the string  $abc$  with regular expression  $r_1$  then

- 1  $Der a (L(r_1))$
- 2  $Der b (Der a (L(r_1)))$
- 3  $Der c (Der b (Der a (L(r_1))))$
- 4 finally we test whether the empty string is in this set; same for  $Der abc (L(r_1))$ .

The matching algorithm works similarly, just over regular expressions instead of sets.



# The Idea with Derivatives

Input: string *abc* and regular expression *r*

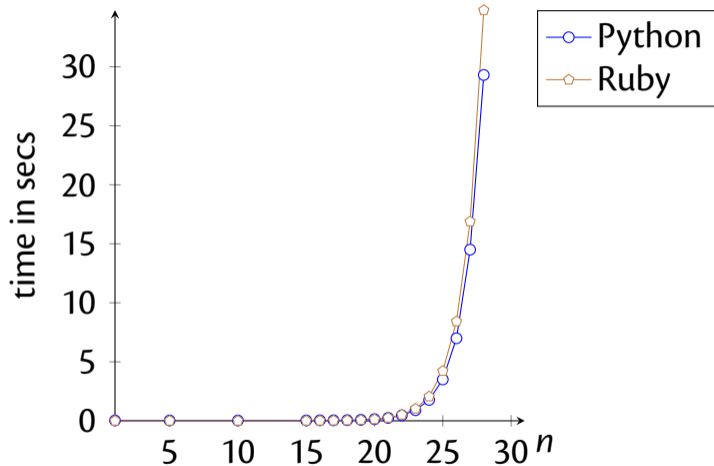
- 1 *der a r*
- 2 *der b (der a r)*
- 3 *der c (der b (der a r))*

# The Idea with Derivatives

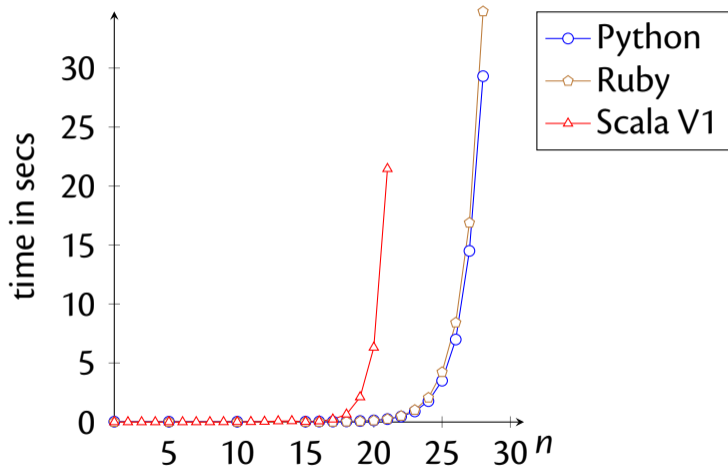
Input: string *abc* and regular expression *r*

- 1 *der a r*
- 2 *der b (der a r)*
- 3 *der c (der b (der a r))*
- 4 finally check whether the last regular expression can match the empty string

$$a^{\{n\}} \cdot a^{\{n\}}$$



# Oops... $a^{\{n\}} \cdot a^{\{n\}}$



# A Problem

We represented the “n-times”  $a^{\{n\}}$  as a sequence regular expression:

0: **1**

1:  $a$

2:  $a \cdot a$

3:  $a \cdot a \cdot a$

...

13:  $a \cdot a \cdot a \cdot a \cdot a \cdot a \cdot a \cdot a \cdot a \cdot a \cdot a \cdot a \cdot a$

...

20:

This problem is aggravated with  $a^?$  being represented as  $a + 1$ .

# Solving the Problem

What happens if we extend our regular expressions with explicit constructors

$$r ::= \dots$$
$$| r^{\{n\}}$$
$$| r^?$$

What is their meaning?

What are the cases for *nullable* and *der*?

# *der* for $n$ -times

Case  $n = 2$  and  $r \cdot r$ :

$$\begin{aligned} \text{der } c(r \cdot r) &\stackrel{\text{def}}{=} \text{if } \text{nullable}(r) \\ &\text{then } (\text{der } c r) \cdot r + \text{der } c r \\ &\text{else } (\text{der } c r) \cdot r \end{aligned}$$

# der for $n$ -times

Case  $n = 2$  and  $r \cdot r$ :

$$\begin{aligned} \text{der } c(r \cdot r) &\stackrel{\text{def}}{=} \text{if } \text{nullable}(r) \\ &\quad \text{then } (\text{der } c r) \cdot r + \text{der } c r \\ &\quad \text{else } (\text{der } c r) \cdot r \end{aligned}$$

$$\begin{aligned} \text{my claim} &\equiv (\text{der } c r) \cdot r \\ \text{(in this case)} & \end{aligned}$$



# der for $n$ -times

Case  $n = 2$  and  $r \cdot r$ :

$$\text{der } c(r \cdot r) \stackrel{\text{def}}{=} \begin{array}{l} \text{if } \text{nullable}(r) \\ \text{then } (\text{der } cr) \cdot r + \text{der } cr \\ \text{else } (\text{der } cr) \cdot r \end{array}$$

my claim  $\equiv (\text{der } cr) \cdot r$   
(in this case)

We know *nullable*(*r*) holds!

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$$(der\ c\ r) \cdot r + der\ c\ r$$

We know *nullable*(*r*) holds!

$$(der\ cr) \cdot r + der\ cr \equiv (der\ cr) \cdot r + (der\ cr) \cdot \mathbf{1}$$

We know *nullable*(*r*) holds!

$$\begin{aligned}(\text{der } c r) \cdot r + \text{der } c r &\equiv (\text{der } c r) \cdot r + (\text{der } c r) \cdot \mathbf{1} \\ &\equiv (\text{der } c r) \cdot (r + \mathbf{1})\end{aligned}$$

We know *nullable*(*r*) holds!

$$\begin{aligned}(\text{der } c r) \cdot r + \text{der } c r &\equiv (\text{der } c r) \cdot r + (\text{der } c r) \cdot \mathbf{1} \\ &\equiv (\text{der } c r) \cdot (r + \mathbf{1}) \\ &\equiv (\text{der } c r) \cdot r \\ &\quad \text{(remember } r \text{ is nullable)}\end{aligned}$$

We know  $\text{nullable}(r)$  holds!

$$\begin{aligned}(\text{der } c r) \cdot r + \text{der } c r &\equiv (\text{der } c r) \cdot r + (\text{der } c r) \cdot \mathbf{1} \\ &\equiv (\text{der } c r) \cdot (r + \mathbf{1}) \\ &\equiv (\text{der } c r) \cdot r \\ &\quad \text{(remember } r \text{ is nullable)}\end{aligned}$$

---

$$\text{der } c (r \cdot r) \stackrel{\text{def}}{=} \begin{array}{l} \text{if } \text{nullable}(r) \\ \text{then } (\text{der } c r) \cdot r + \text{der } c r \\ \text{else } (\text{der } c r) \cdot r \end{array}$$

We know *nullable*(*r*) holds!

$$\begin{aligned}(\text{der } c \ r) \cdot r + \text{der } c \ r &\equiv (\text{der } c \ r) \cdot r + (\text{der } c \ r) \cdot \mathbf{1} \\ &\equiv (\text{der } c \ r) \cdot (r + \mathbf{1}) \\ &\equiv (\text{der } c \ r) \cdot r \\ &\quad \text{(remember } r \text{ is nullable)}\end{aligned}$$

---

$$\text{der } c \ (r \cdot r) \stackrel{\text{def}}{=} \begin{array}{l} \text{if } \text{nullable}(r) \\ \text{then } (\text{der } c \ r) \cdot r \\ \text{else } (\text{der } c \ r) \cdot r \end{array}$$



We know *nullable*(*r*) holds!

$$\begin{aligned}(\text{der } c r) \cdot r + \text{der } c r &\equiv (\text{der } c r) \cdot r + (\text{der } c r) \cdot \mathbf{1} \\ &\equiv (\text{der } c r) \cdot (r + \mathbf{1}) \\ &\equiv (\text{der } c r) \cdot r \\ &\quad \text{(remember } r \text{ is nullable)}\end{aligned}$$

---

$$\text{der } c (r \cdot r) \stackrel{\text{def}}{=} (\text{der } c r) \cdot r$$

	$r\{n\}$	$der$
$n = 0:$	<b>1</b>	<b>0</b>
$n = 1:$	$r$	$(der\ r)$
$n = 2:$	$r \cdot r$	$(der\ r) \cdot r$
$n = 3:$	$r \cdot r \cdot r$	???
	$\vdots$	

	$r\{n\}$	$der$
$n = 0:$	<b>1</b>	<b>0</b>
$n = 1:$	$r$	$(der\ c\ r)$
$n = 2:$	$r \cdot r$	$(der\ c\ r) \cdot r$
$n = 3:$	$r \cdot r \cdot r$	$(der\ c\ r) \cdot r \cdot r$
	$\vdots$	

	$r\{n\}$	$der$
$n = 0:$	<b>1</b>	<b>0</b>
$n = 1:$	$r$	$(der\ c\ r)$
$n = 2:$	$r \cdot r$	$(der\ c\ r) \cdot r$
$n = 3:$	$r \cdot r \cdot r$	$(der\ c\ r) \cdot r \cdot r$
	$\vdots$	

$nullable(r\{n\}) \stackrel{\text{def}}{=} \text{if } n = 0 \text{ then } true \text{ else } nullable(r)$

$der\ c\ (r\{n\}) \stackrel{\text{def}}{=} \text{if } n = 0 \text{ then } \mathbf{0} \text{ else } (der\ c\ r) \cdot r\{n - 1\}$

$derc(r \cdot r \cdot r) \stackrel{\text{def}}{=} \begin{cases} \text{if } nullable(r) \\ \text{then } (dercr) \cdot r \cdot r + derc(r \cdot r) \\ \text{else } (dercr) \cdot r \cdot r \end{cases}$

$derc(r \cdot r \cdot r) \stackrel{\text{def}}{=} \begin{array}{l} \text{if } nullable(r) \\ \text{then } (derc\ r) \cdot r \cdot r + (derc\ r) \cdot r \\ \text{else } (derc\ r) \cdot r \cdot r \end{array}$

$derc(r \cdot r \cdot r) \stackrel{\text{def}}{=} \begin{cases} \text{if } nullable(r) \\ \text{then } (derc\ r) \cdot (r \cdot r + r) \\ \text{else } (derc\ r) \cdot r \cdot r \end{cases}$

$derc(r \cdot r \cdot r) \stackrel{\text{def}}{=} \begin{array}{l} \text{if } nullable(r) \\ \text{then } (derc\ r) \cdot (r \cdot (r + \mathbf{1})) \\ \text{else } (derc\ r) \cdot r \cdot r \end{array}$

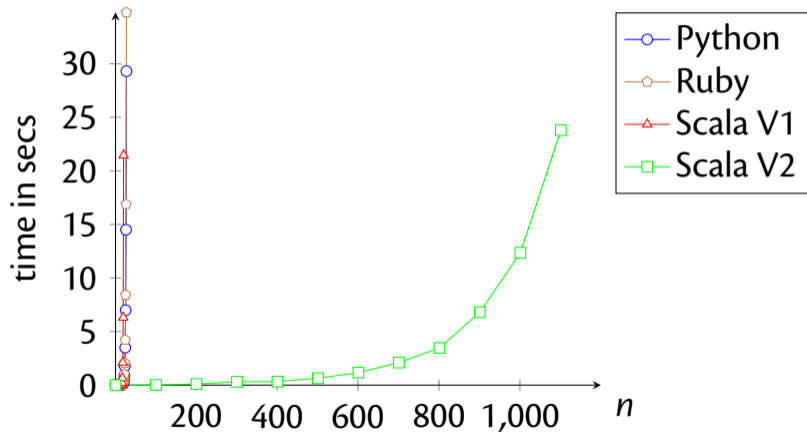


$derc(r \cdot r \cdot r) \stackrel{\text{def}}{=} \begin{array}{l} \text{if } nullable(r) \\ \text{then } (derc\ r) \cdot (r \cdot r) \\ \text{else } (derc\ r) \cdot r \cdot r \end{array}$

$derc(r \cdot r \cdot r) \stackrel{\text{def}}{=} \begin{array}{l} \text{if } nullable(r) \\ \text{then } (derc\ r) \cdot r \cdot r \\ \text{else } (derc\ r) \cdot r \cdot r \end{array}$

$$\text{der } c(r \cdot r \cdot r) \stackrel{\text{def}}{=} (\text{der } cr) \cdot r \cdot r$$

# Brzozowski: $a^? \{n\} \cdot a \{n\}$



# Examples

Recall the example of  $r \stackrel{\text{def}}{=} ((a \cdot b) + b)^*$  with

$$\text{der } a r = ((\mathbf{1} \cdot b) + \mathbf{0}) \cdot r$$

$$\text{der } b r = ((\mathbf{0} \cdot b) + \mathbf{1}) \cdot r$$

$$\text{der } c r = ((\mathbf{0} \cdot b) + \mathbf{0}) \cdot r$$

What are these regular expressions equivalent to?

# Simplification Rules

$$r + 0 \Rightarrow r$$

$$0 + r \Rightarrow r$$

$$r \cdot 1 \Rightarrow r$$

$$1 \cdot r \Rightarrow r$$

$$r \cdot 0 \Rightarrow 0$$

$$0 \cdot r \Rightarrow 0$$

$$r + r \Rightarrow r$$

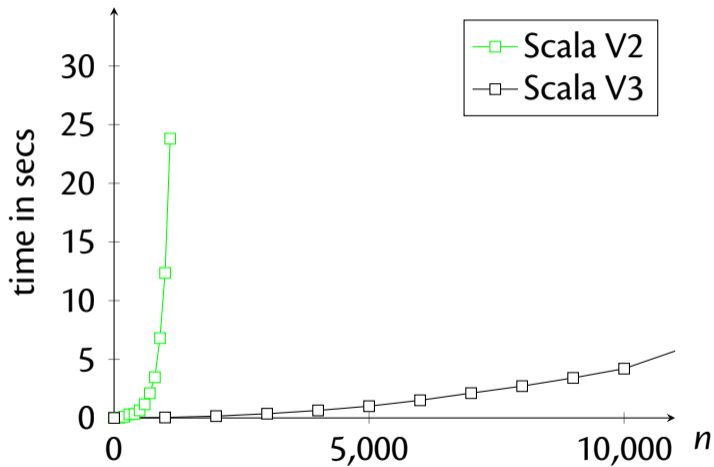
```
def ders(s: List[Char], r: Rexp) : Rexp = s match {  
  case Nil => r  
  case c::s => ders(s, simp(der(c, r)))  
}
```

```

def simp(r: Rexp) : Rexp = r match {
  case ALT(r1, r2) => {
    (simp(r1), simp(r2)) match {
      case (ZERO, r2s) => r2s
      case (r1s, ZERO) => r1s
      case (r1s, r2s) =>
        if (r1s == r2s) r1s else ALT(r1s, r2s)
    }
  }
  case SEQ(r1, r2) => {
    (simp(r1), simp(r2)) match {
      case (ZERO, _) => ZERO
      case (_, ZERO) => ZERO
      case (ONE, r2s) => r2s
      case (r1s, ONE) => r1s
      case (r1s, r2s) => SEQ(r1s, r2s)
    }
  }
  case r => r
}

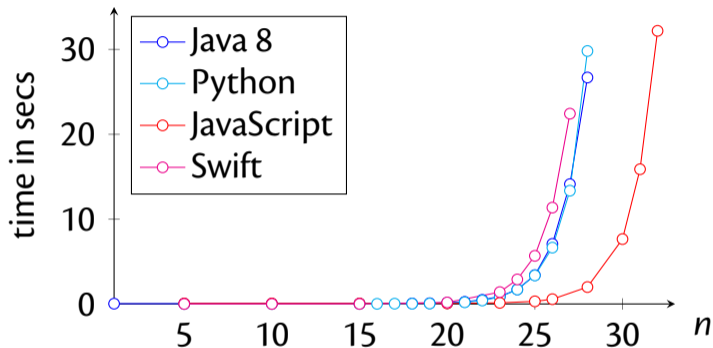
```

# Brzozowski: $a^? \{n\} \cdot a \{n\}$





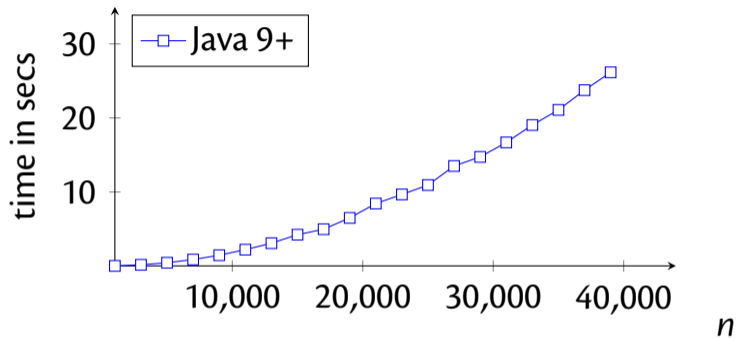
# Another Example $(a^*)^* \cdot b$



Regex:  $(a^*)^* \cdot b$

Strings of the form  $\underbrace{a \dots a}_n$

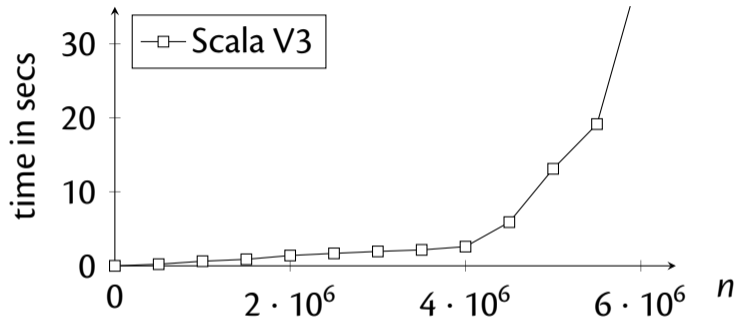
# Same Example in Java 9+



Regex:  $(a^*)^* \cdot b$

Strings of the form  $\underbrace{a \dots a}_n$

## ...and with Brzowski



Regex:  $(a^*)^* \cdot b$

Strings of the form  $\underbrace{a \dots a}_n$

# What is good about this Alg.

- extends to most regular expressions, for example  $\sim r$  (next slide)
- is easy to implement in a functional language
- the algorithm is already quite old; there is still work to be done to use it as a tokenizer (that is relatively new work)
- we can prove its correctness...(another video)

# Negation of Regular Expr's

- $\sim r$  (everything that  $r$  cannot recognise)
- $L(\sim r) \stackrel{\text{def}}{=} UNIV - L(r)$
- $nullable(\sim r) \stackrel{\text{def}}{=} \text{not}(nullable(r))$
- $derc(\sim r) \stackrel{\text{def}}{=} \sim(derc r)$

# Negation of Regular Expr's

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- $derc(\sim r) \stackrel{\text{def}}{=} \sim(derc r)$

Used often for recognising comments:

$$/ \cdot * \cdot (\sim ([a-z]^* \cdot * \cdot / \cdot [a-z]^*)) \cdot * \cdot /$$

# Coursework 1

- Submission on Friday 16 October @ 18:00
- source code needs to be submitted as well
- you can re-use my Scala code from KEATS and use any programming language you like
- <https://nms.kcl.ac.uk/christian.urban/ProgInScala2ed.pdf>

# Proofs about Rexprs

Remember their inductive definition:

$$r ::= \begin{array}{l} 0 \\ 1 \\ c \\ r_1 \cdot r_2 \\ r_1 + r_2 \\ r^* \end{array}$$

If we want to prove something, say a property  $P(r)$ ,  
for all regular expressions  $r$  then ...



# Proofs about Rexp (2)

- $P$  holds for  $0$ ,  $1$  and  $c$
- $P$  holds for  $r_1 + r_2$  under the assumption that  $P$  already holds for  $r_1$  and  $r_2$ .
- $P$  holds for  $r_1 \cdot r_2$  under the assumption that  $P$  already holds for  $r_1$  and  $r_2$ .
- $P$  holds for  $r^*$  under the assumption that  $P$  already holds for  $r$ .

# Proofs about Rexp (3)

Assume  $P(r)$  is the property:

$nullable(r)$  if and only if  $[\ ] \in L(r)$

# Proofs about Rexp (4)

$$\text{rev}(\mathbf{0}) \stackrel{\text{def}}{=} \mathbf{0}$$

$$\text{rev}(\mathbf{1}) \stackrel{\text{def}}{=} \mathbf{1}$$

$$\text{rev}(c) \stackrel{\text{def}}{=} c$$

$$\text{rev}(r_1 + r_2) \stackrel{\text{def}}{=} \text{rev}(r_1) + \text{rev}(r_2)$$

$$\text{rev}(r_1 \cdot r_2) \stackrel{\text{def}}{=} \text{rev}(r_2) \cdot \text{rev}(r_1)$$

$$\text{rev}(r^*) \stackrel{\text{def}}{=} \text{rev}(r)^*$$

We can prove

$$L(\text{rev}(r)) = \{s^{-1} \mid s \in L(r)\}$$

by induction on  $r$ .

# Correctness Proof for our Matcher

- We started from

$$s \in L(r)$$

$$\Leftrightarrow \square \in \text{Ders } s (L(r))$$

# Correctness Proof for our Matcher

- We started from

$$s \in L(r)$$

$$\Leftrightarrow [] \in \text{Ders } s (L(r))$$

- if we can show  $\text{Ders } s (L(r)) = L(\text{ders } s r)$  we have

$$\Leftrightarrow [] \in L(\text{ders } s r)$$

$$\Leftrightarrow \text{nullable}(\text{ders } s r)$$

$$\stackrel{\text{def}}{=} \text{matcher } s r$$

# Proofs about Rexp (5)

Let  $Der\ c\ A$  be the set defined as

$$Der\ c\ A \stackrel{\text{def}}{=} \{s \mid c :: s \in A\}$$

We can prove

$$L(\text{der}\ c\ r) = Der\ c\ (L(r))$$

by induction on  $r$ .

# Proofs about Strings

If we want to prove something, say a property  $P(s)$ , for all strings  $s$  then ...

- $P$  holds for the empty string, and
- $P$  holds for the string  $c::s$  under the assumption that  $P$  already holds for  $s$

# Proofs about Strings (2)

We can then prove

$$\text{Ders } s (L(r)) = L(\text{ders } s r)$$

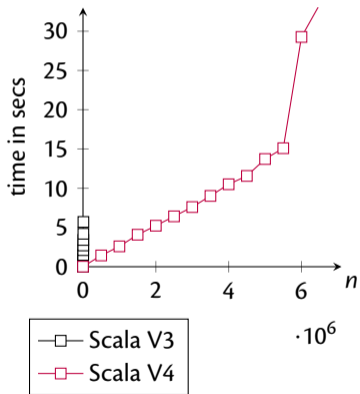
We can finally prove

$$\text{matcher } s r \text{ if and only if } s \in L(r)$$

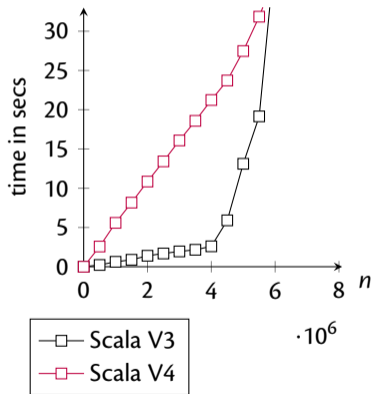


# Epilogue

Graph:  $a^{a^n} \cdot a^n$



Graph:  $(a^*)^* \cdot b$

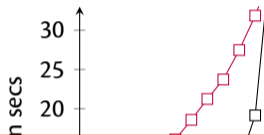


# Epilogue

Graph:  $a^?{n} \cdot a^{n}$



Graph:  $(a^*)^* \cdot b$



```
def ders2(s: List[Char], r: Rexp) : Rexp = (s, r) match {  
  case (Nil, r) => r  
  case (s, ZERO) => ZERO  
  case (s, ONE) => if (s == Nil) ONE else ZERO  
  case (s, CHAR(c)) => if (s == List(c)) ONE else  
    if (s == Nil) CHAR(c) else ZERO  
  case (s, ALT(r1, r2)) => ALT(ders2(s, r1), ders2(s, r2))  
  case (c::s, r) => ders2(s, simp(der(c, r)))  
}
```

# Another Homework Question

- How many basic regular expressions are there to match the string *abcd*?

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- How many if they are also not allowed to contain stars?
- How many if they are also not allowed to contain  $\_ + \_$ ?

**Questions?**

# Last Week

Last week I showed you a regular expression matcher that works provably correct in all cases (we only started with the proving part though)

*matcher*  $s$   $r$  if and only if  $s \in L(r)$

by Janusz Brzozowski (1964)



# Proofs about Rexp

- $P$  holds for  $0$ ,  $1$  and  $c$
- $P$  holds for  $r_1 + r_2$  under the assumption that  $P$  already holds for  $r_1$  and  $r_2$ .
- $P$  holds for  $r_1 \cdot r_2$  under the assumption that  $P$  already holds for  $r_1$  and  $r_2$ .
- $P$  holds for  $r^*$  under the assumption that  $P$  already holds for  $r$ .

We proved

*nullable*(*r*) if and only if  $\epsilon \in L(r)$

by induction on the regular expression *r*.

We proved

*nullable*( $r$ ) if and only if  $\epsilon \in L(r)$

by induction on the regular expression  $r$ .

# Any Questions?

# Proofs about Natural Numbers and Strings

- $P$  holds for  $0$  and
- $P$  holds for  $n + 1$  under the assumption that  $P$  already holds for  $n$
  
- $P$  holds for  $[]$  and
- $P$  holds for  $c :: s$  under the assumption that  $P$  already holds for  $s$

# Correctness Proof for our Matcher

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# Correctness Proof for our Matcher

- We started from

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- **if** we can show  $\text{Ders } s (L(r)) = L(\text{ders } s r)$  we have

$$\Leftrightarrow [] \in L(\text{ders } s r)$$

$$\Leftrightarrow \text{nullable}(\text{ders } s r)$$

$$\stackrel{\text{def}}{=} \text{matcher } s r$$

We need to prove

$$L(\text{der } c r) = \text{Der } c (L(r))$$

also by induction on the regular expression  $r$ .