Homework 7

Please submit your solutions via email. Please submit only ASCII text or PDFs. Every solution should be preceeded by the corresponding question, like:

Qn:...a difficult question from me...A:...an answer from you ...Qn+1...another difficult question...A:...another brilliant answer from you...

Solutions will only be accepted until 30th December!

1. Suppose the context-sensitive grammar

$$\begin{array}{cccc} S &
ightarrow & bSAA \mid \epsilon \ A &
ightarrow & a \ bA &
ightarrow & Ab \end{array}$$

where *S* is the starting symbol of the grammar. Give a derivation of the string *"aaabaaabb"*. What can you say about the number of as and bs in the strings recognised by this grammar.

2. Consider the following grammar

$$\begin{array}{l} S \rightarrow N \cdot P \\ P \rightarrow V \cdot N \\ N \rightarrow N \cdot N \\ N \rightarrow A \cdot N \\ N \rightarrow \text{student} \mid \text{trainer} \mid \text{team} \mid \text{trains} \\ V \rightarrow \text{trains} \mid \text{team} \\ A \rightarrow \text{The} \mid \text{the} \end{array}$$

where *S* is the start symbol and *S*, *P*, *N*, *V* and *A* are non-terminals. Using the CYK-algorithm, check whether or not the following string can be parsed by the grammar:

The trainer trains the student team

3. Transform the grammar

$$\begin{array}{rrrr} A & \rightarrow & 0A1 \mid BB \\ B & \rightarrow & \epsilon \mid 2B \end{array}$$

into Chomsky normal form.

4. Consider the following grammar *G*

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\begin{array}{l} S \rightarrow \texttt{if0} \cdot E \cdot \texttt{then} \cdot S \\ S \rightarrow \texttt{print} \cdot S \\ S \rightarrow \texttt{begin} \cdot B \cdot \texttt{end} \\ B \rightarrow S \cdot \texttt{;} \\ B \rightarrow S \cdot \texttt{;} \cdot B \\ E \rightarrow num \end{array}
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where *S* is the start symbol and *S*, *E* and *B* are non-terminals.

Check each rule below and decide whether, when added to *G*, the combined grammar is ambiguous. If yes, give a string that has more than one parse tree.

 $\begin{array}{ll} (\mathrm{i}) & S \rightarrow \mathtt{if0} \cdot E \cdot \mathtt{then} \cdot S \cdot \mathtt{else} \cdot S \\ (\mathrm{ii}) & B \rightarrow B \cdot B \\ (\mathrm{iii}) & E \rightarrow (\cdot E \cdot) \\ (\mathrm{iv}) & E \rightarrow E \cdot + \cdot E \end{array}$