Automata and Formal Languages (3)

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Slides: KEATS (also home work and course-work is there)

Regular Expressions

In programming languages they are often used to recognise:

- symbols, digits
- identifiers
- numbers (non-leading zeros)
- keywords
- comments

`http://www.regexper.com`

Last Week

Last week I showed you a regular expression matcher which works provably correct in all cases (we did not do the proving part though)

$$\mathit{matches}\, r\, s \quad \text{if and only if} \quad s \in L(r)$$

by Janusz Brzozowski (1964)

The Derivative of a Rexp

$$
\begin{array}{lcl}
der\, c\, (\varnothing) & \stackrel{\text{def}}{=} & \varnothing \\
der\, c\, (\epsilon) & \stackrel{\text{def}}{=} & \varnothing \\
der\, c\, (d) & \stackrel{\text{def}}{=} & \text{if } c = d \text{ then } \epsilon \text{ else } \varnothing \\
der\, c\, (r_1 + r_2) & \stackrel{\text{def}}{=} & der\, c\, r_1 + der\, c\, r_2 \\
der\, c\, (r_1 \cdot r_2) & \stackrel{\text{def}}{=} & \text{if } nullable(r_1) \\
 & & \text{then } (der\, c\, r_1) \cdot r_2 + der\, c\, r_2 \\
 & & \text{else } (der\, c\, r_1) \cdot r_2 \\
der\, c\, (r^*) & \stackrel{\text{def}}{=} & (der\, c\, r) \cdot (r^*) \\
 & & \\
ders\, []\, r & \stackrel{\text{def}}{=} & r \\
ders\, (c :: s)\, r & \stackrel{\text{def}}{=} & ders\, s\, (der\, c\, r)
\end{array}
$$

To see what is going on, define

$$Der\, c\, A \stackrel{\text{def}}{=} \{s \mid c :: s \in A\}$$

For $A = \{foo, bar, frak\}$ then

$$\begin{array}{lcl} Der\, f\, A & = & \{oo, rak\} \\ Der\, b\, A & = & \{ar\} \\ Der\, a\, A & = & \varnothing \end{array}$$

The Idea of the Algorithm

If we want to recognise the string abc with regular expression r then

1. $Der\, a\, (L(r))$

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3. $Der\, c\, (Der\, b\, (Der\, a\, (L(r))))$

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The matching algorithm works similarly, just over regular expressions instead of sets.

Input: string abc and regular expression r

1. $der\, a\, r$
2. $der\, b\, (der\, a\, r)$
3. $der\, c\, (der\, b\, (der\, a\, r))$

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1. $der\, a\, r$
2. $der\, b\, (der\, a\, r)$
3. $der\, c\, (der\, b\, (der\, a\, r))$

4. finally check whether the last regular expression can match the empty string

We proved already

$$\mathit{nullable}(r) \text{ if and only if } [] \in L(r)$$

by induction on the regular expression.

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Any Questions?

We need to prove

$$L(der\, c\, r) = Der\, c\, (L(r))$$

by induction on the regular expression.

Proofs about Rexps

- P holds for $\varnothing$, ϵ and c
- P holds for $r_1 + r_2$ under the assumption that P already holds for r_1 and r_2.
- P holds for $r_1 \cdot r_2$ under the assumption that P already holds for r_1 and r_2.
- P holds for r^* under the assumption that P already holds for r.

Proofs about Natural Numbers and Strings

- P holds for 0 and
- P holds for $n + 1$ under the assumption that P already holds for n

- P holds for $[]$ and
- P holds for $c :: s$ under the assumption that P already holds for s

Languages

A language is a set of strings.

A regular expression specifies a language.

A language is regular iff there exists a regular expression that recognises all its strings.

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A language is regular iff there exists a regular expression that recognises all its strings.

not all languages are regular, e.g. $a^n b^n$ is not

Regular Expressions

$$
\begin{array}{lcll}
r & ::= & \varnothing & \text{null} \\
 & | & \epsilon & \text{empty string / "" / []} \\
 & | & c & \text{character} \\
 & | & r_1 \cdot r_2 & \text{sequence} \\
 & | & r_1 + r_2 & \text{alternative / choice} \\
 & | & r^* & \text{star (zero or more)}
\end{array}
$$

How about ranges $[a\text{-}z]$, r^+ and $\sim r$? Do they increase the set of languages we can recognise?

Negation of Regular Expr's

- $\sim r$ (everything that r cannot recognise)
- $L(\sim r) \stackrel{\text{def}}{=} UNIV - L(r)$
- $nullable(\sim r) \stackrel{\text{def}}{=} not\,(nullable(r))$
- $der\,c\,(\sim r) \stackrel{\text{def}}{=} \sim (der\,c\,r)$

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Used often for recognising comments:

$$/ \cdot * \cdot (\sim ([a\text{-}z]^* \cdot * \cdot / \cdot [a\text{-}z]^*)) \cdot * \cdot /$$

Negation

Assume you have an alphabet consisting of the letters *a*, *b* and *c* only. Find a (basic!) regular expression that matches all strings *except ab* and *ac*!

Automata

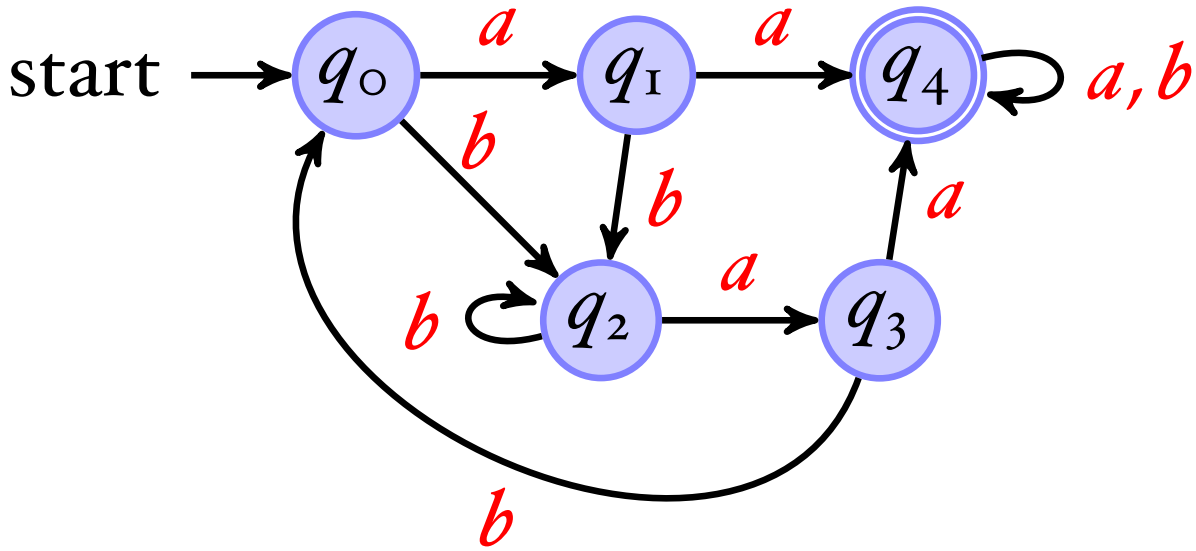
A deterministic finite automaton consists of:

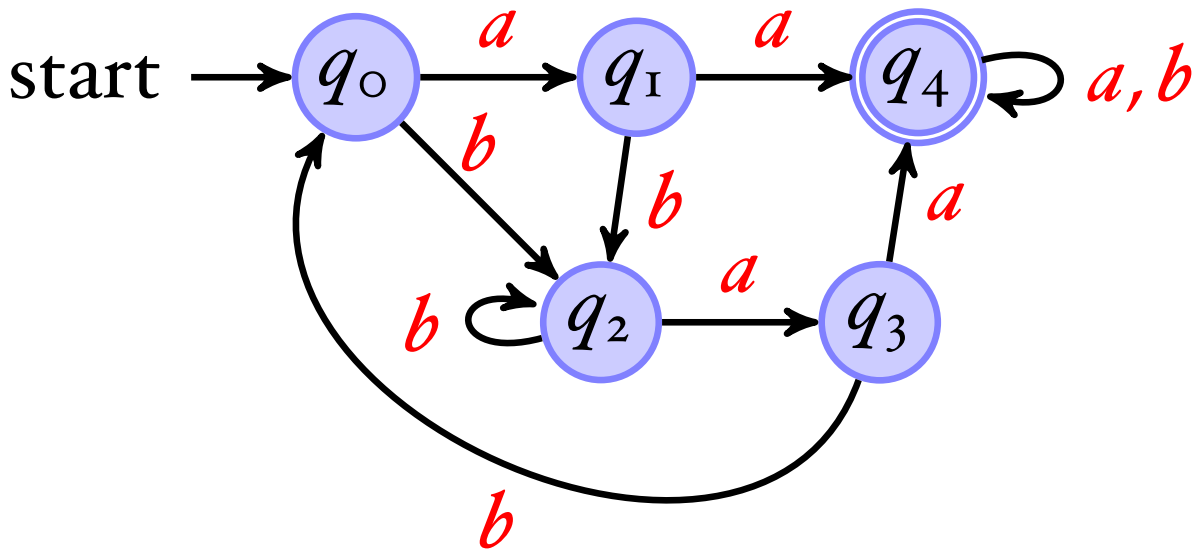
- a set of states
- one of these states is the start state
- some states are accepting states, and
- there is transition function

  which takes a state as argument and a character and produces a new state

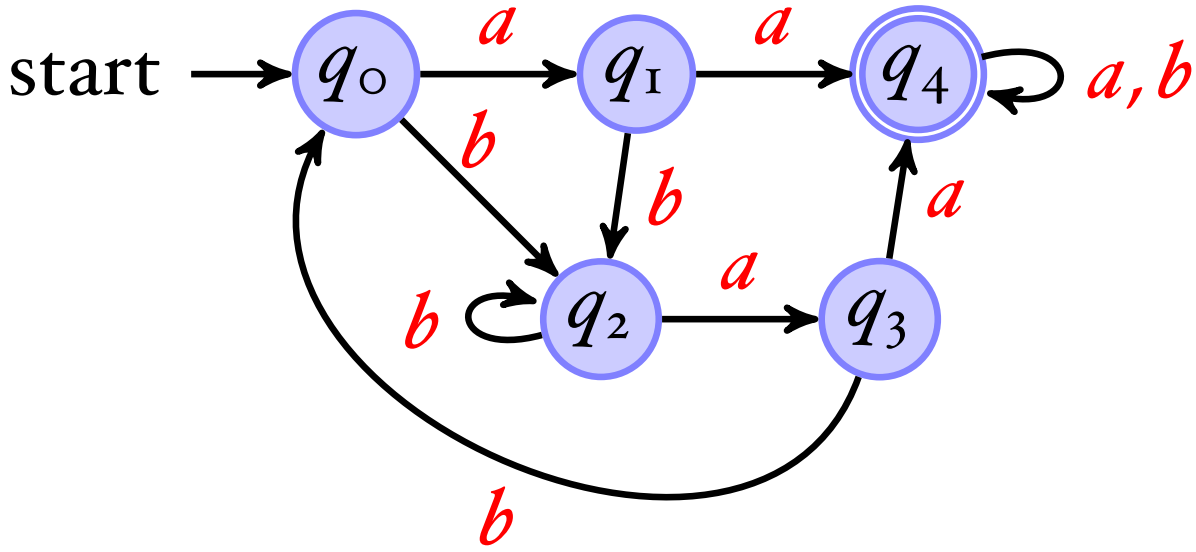
  this function might not be everywhere defined

$$A(\mathcal{Q}, q_0, F, \delta)$$


start
q_0
q_1
q_4
q_2
q_3
a
a
a, b
b
b
a
b
a
b



- the start state can be an accepting state
- it is possible that there is no accepting state
- all states might be accepting (but this does not necessarily mean all strings are accepted)



for this automaton δ is the function

$$\begin{array}{lll} (q_0, a) \to q_1 & (q_1, a) \to q_4 & (q_4, a) \to q_4 \\ (q_0, b) \to q_2 & (q_1, b) \to q_2 & (q_4, b) \to q_4 \end{array} \quad \ldots$$

Accepting a String

Given

$$A(\mathcal{Q}, q_0, F, \delta)$$

you can define

$$\hat{\delta}(q, []) \stackrel{\text{def}}{=} q$$
$$\hat{\delta}(q, c :: s) \stackrel{\text{def}}{=} \hat{\delta}(\delta(q, c), s)$$

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Whether a string s is accepted by A?

$$\hat{\delta}(q_0, s) \in F$$

Non-Deterministic Finite Automata

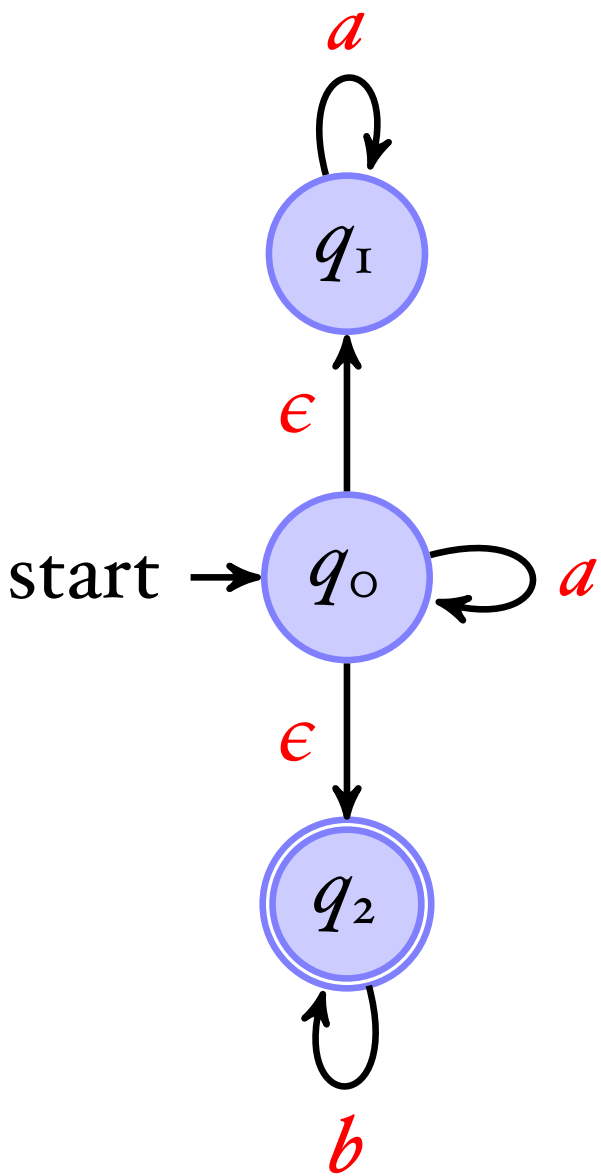
A non-deterministic finite automaton consists again of:

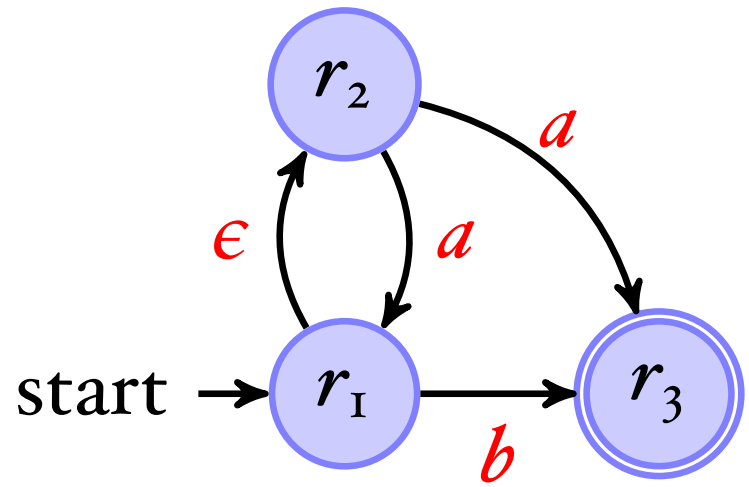
- a finite set of states
- one of these states is the start state
- some states are accepting states, and
- there is transition relation

$$(q_1, a) \rightarrow q_2$$
$$(q_1, a) \rightarrow q_3$$

$$(q_1, \epsilon) \rightarrow q_2$$

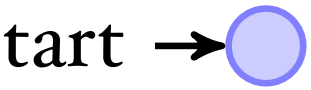
Two NFA Examples



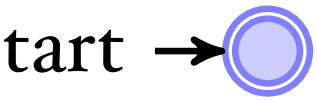


Rexp to NFA

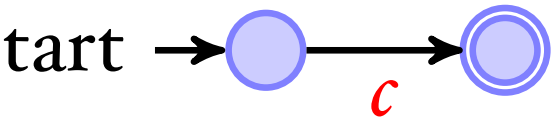
$\varnothing$



ϵ

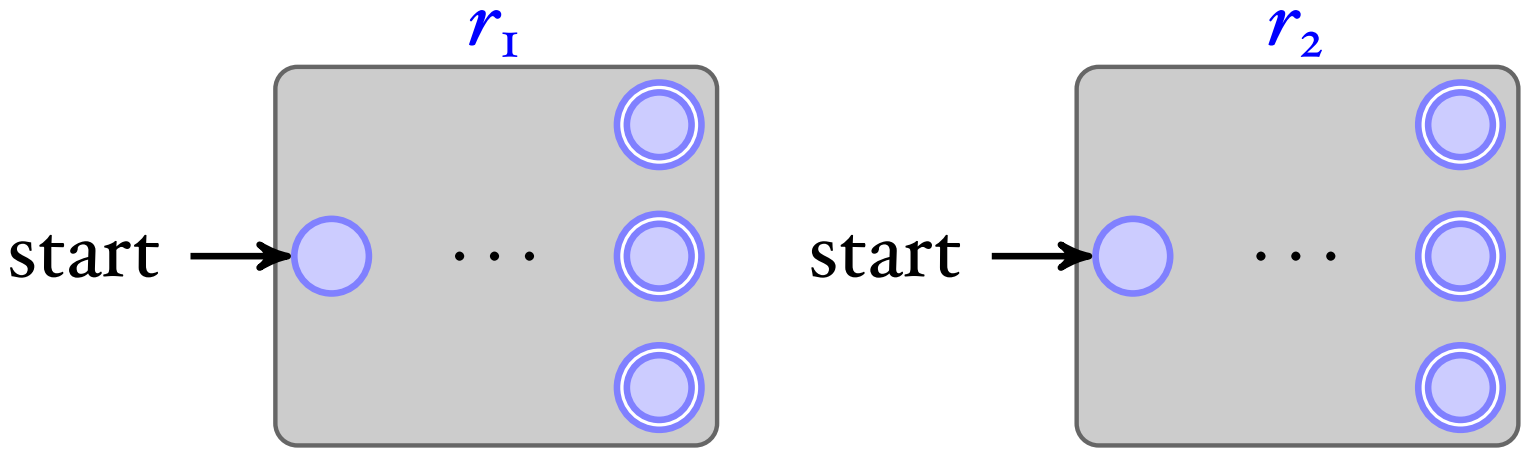


c



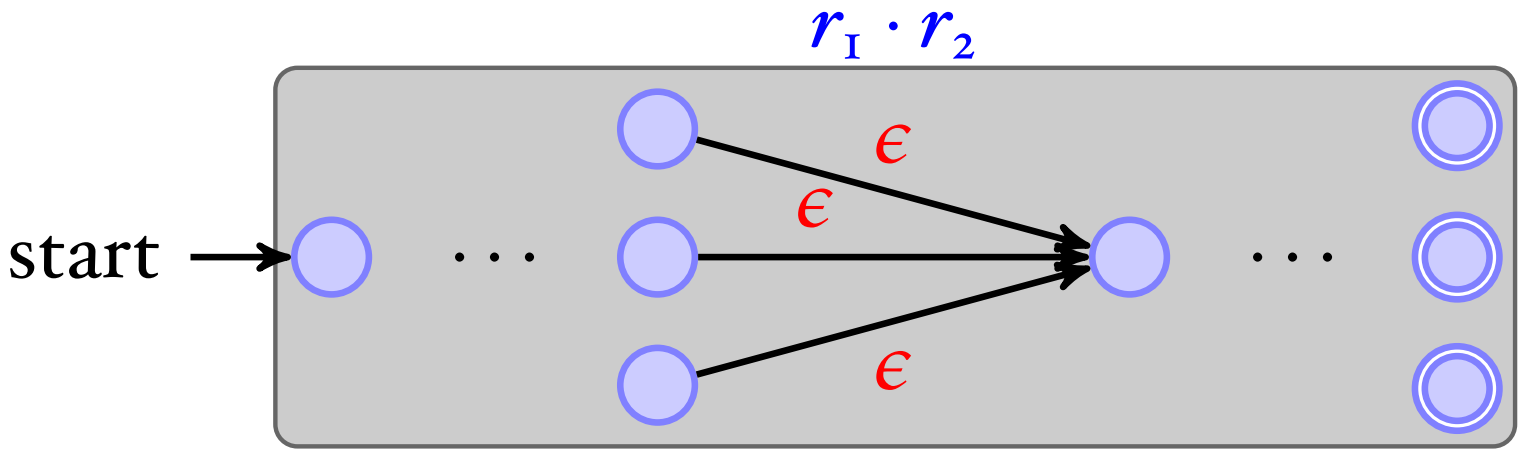
Case $r_1 \cdot r_2$

By recursion we are given two automata:



We need to (1) change the accepting nodes of the first automaton into non-accepting nodes, and (2) connect them via ϵ-transitions to the starting state of the second automaton.

Case $r_1 \cdot r_2$



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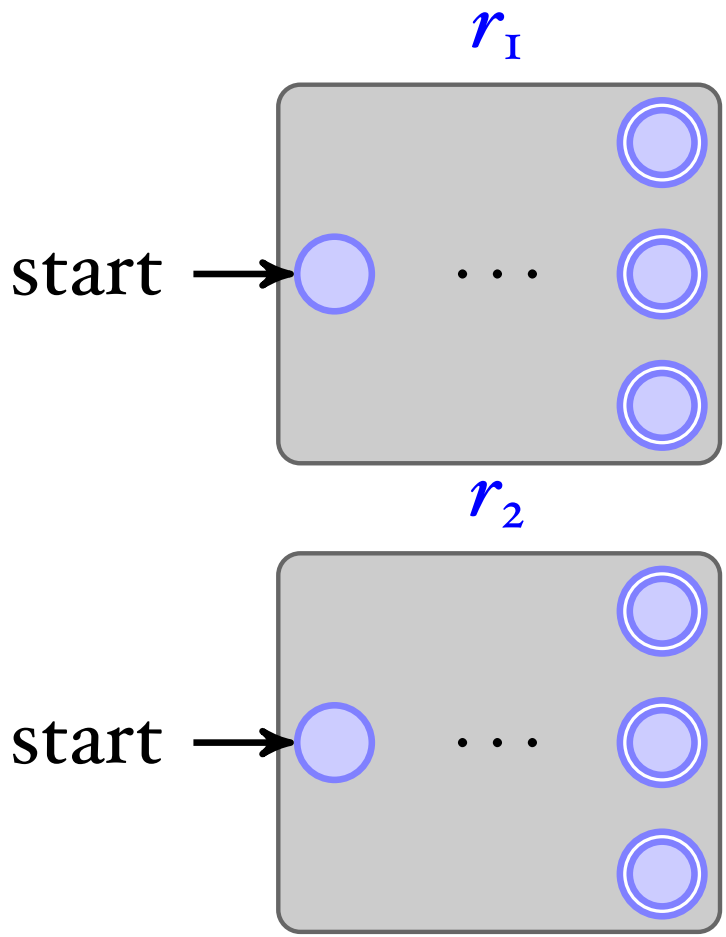

We (1) need to introduce a new starting state and (2) connect it to the original two starting states.

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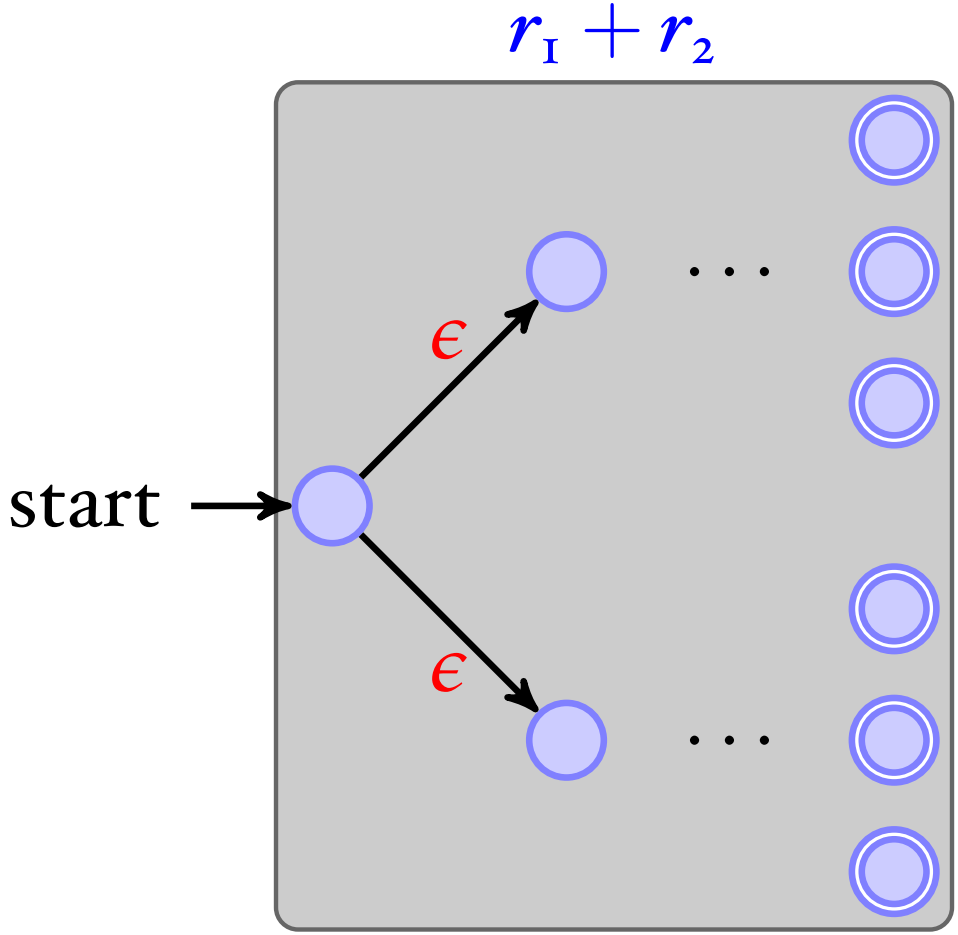
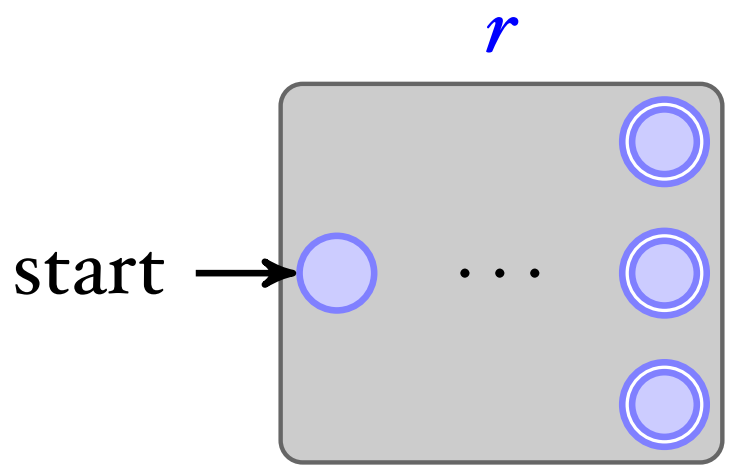


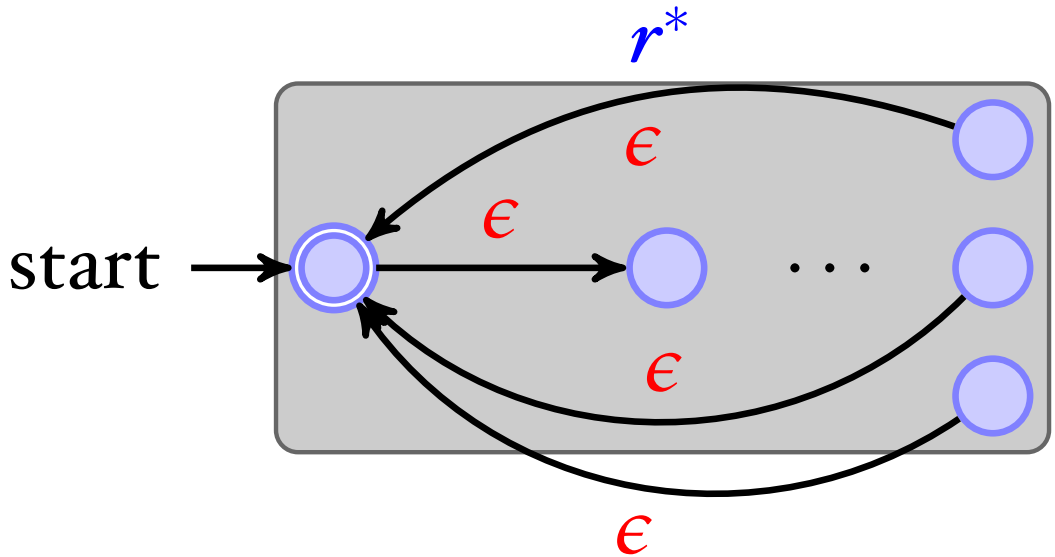
We (1) need to introduce a new starting state and (2) connect it to the original two starting states.

Case r^*

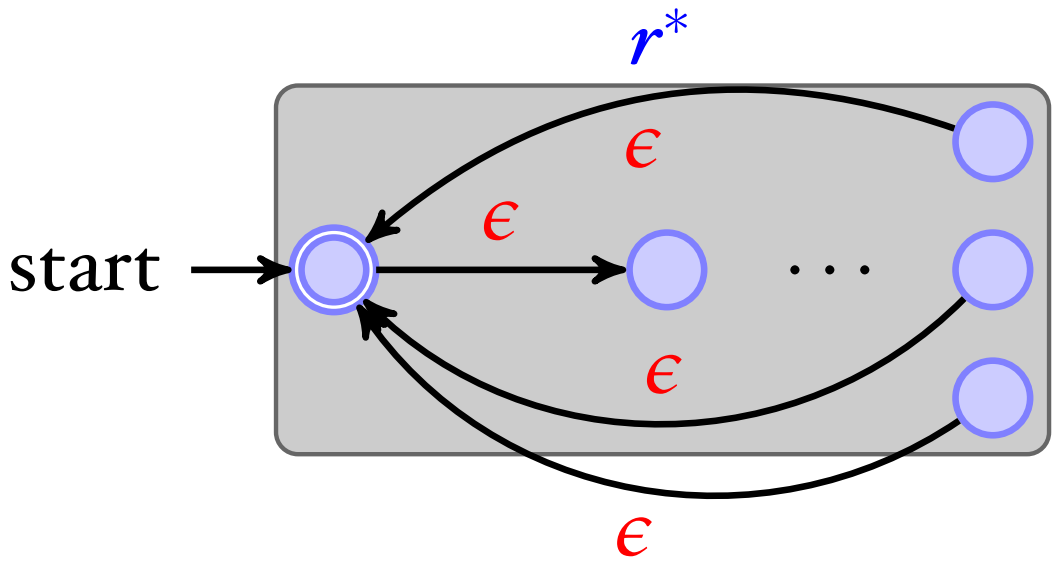
By recursion we are given an automaton for r:



Case r^*

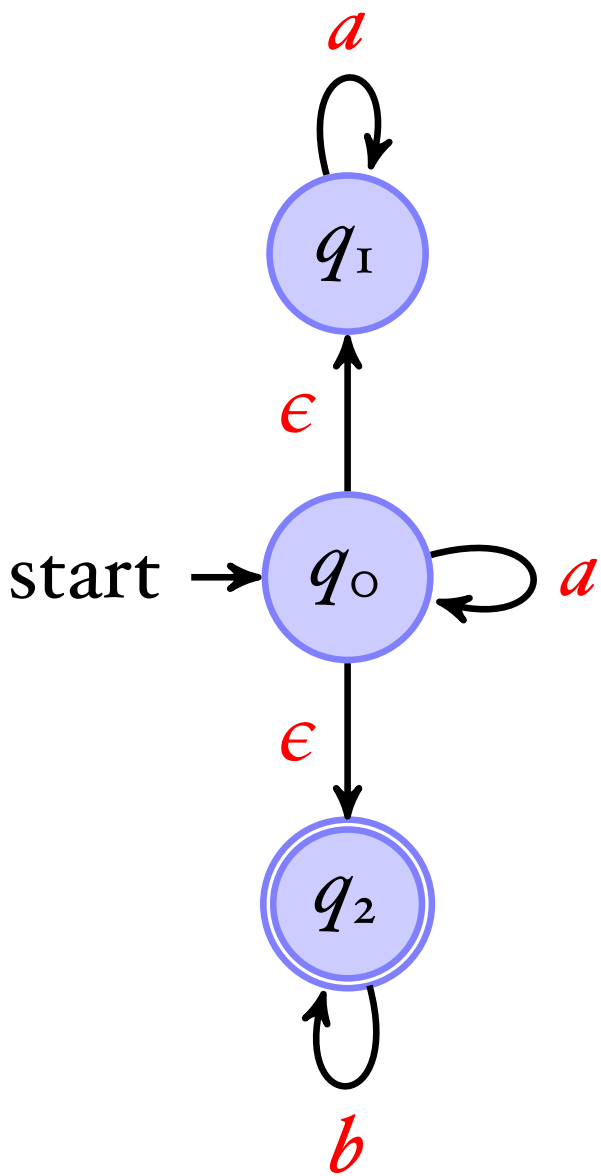


Case r^*



Why can't we just have an epsilon transition from the accepting states to the starting state?

Subset Construction



nodes	a	b
$\emptyset$		
$\{0\}$		
$\{1\}$		
$\{2\}$		
$\{0,1\}$		
$\{0,2\}$		
$\{1,2\}$		
$\{0,1,2\}$		

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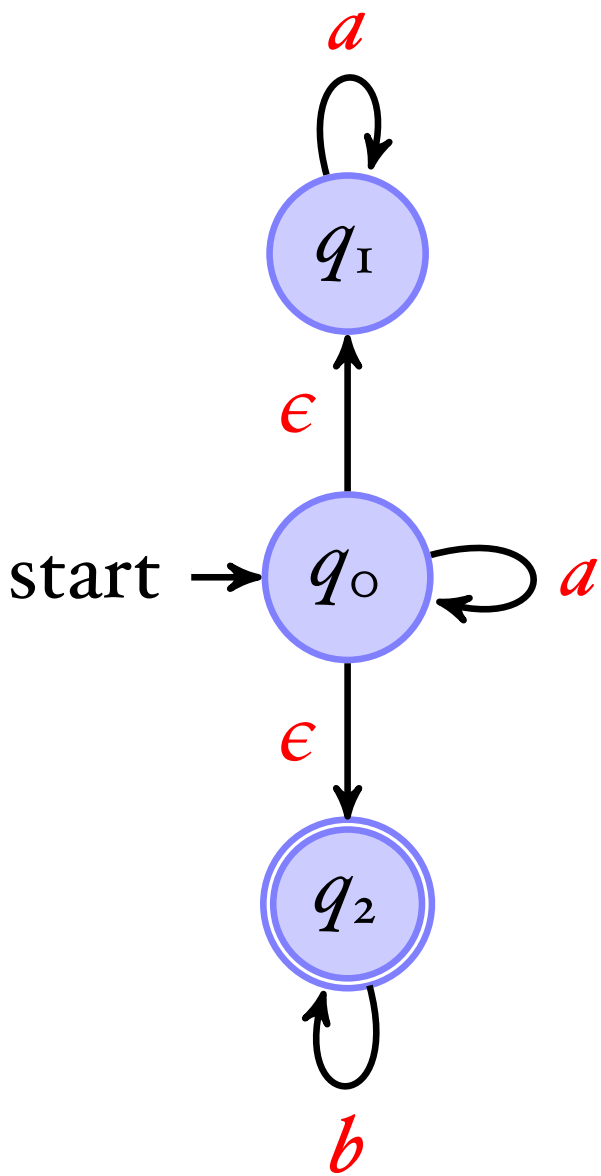

nodes	a	b
$\emptyset$	$\emptyset$	$\emptyset$
$\{0\}$		
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$\{2\}$		
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$\{1,2\}$		
$\{0,1,2\}$		

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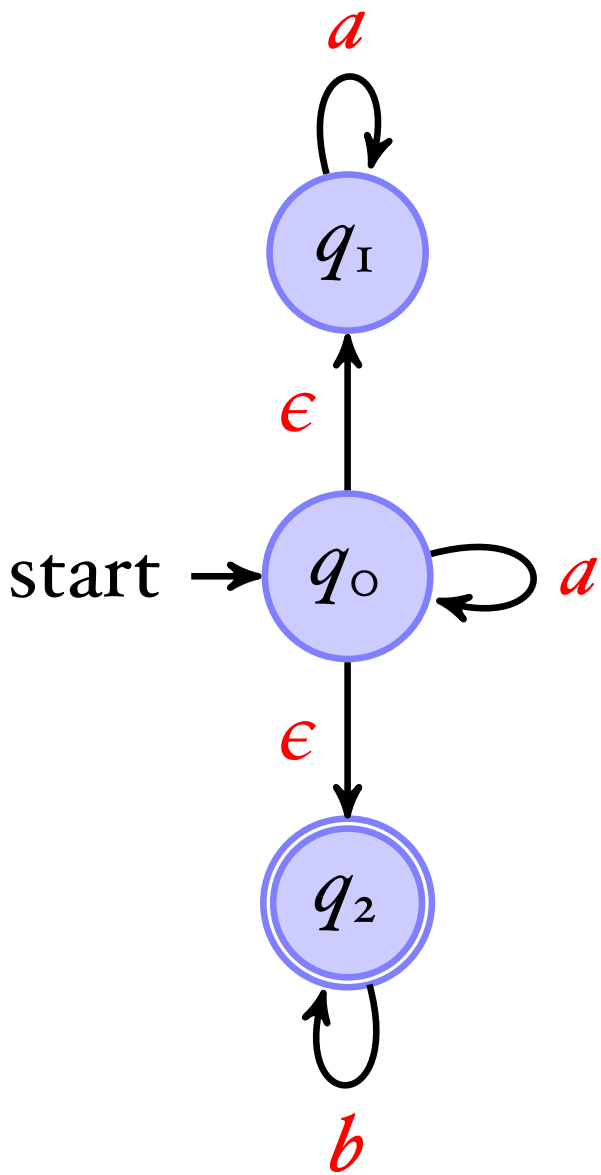

nodes	a	b
$\varnothing$	$\varnothing$	$\varnothing$
$\{0\}$	$\{0,1,2\}$	$\{2\}$
$\{1\}$	$\{1\}$	$\varnothing$
$\{2\}$	$\varnothing$	$\{2\}$
$\{0,1\}$		
$\{0,2\}$		
$\{1,2\}$		
$\{0,1,2\}$		

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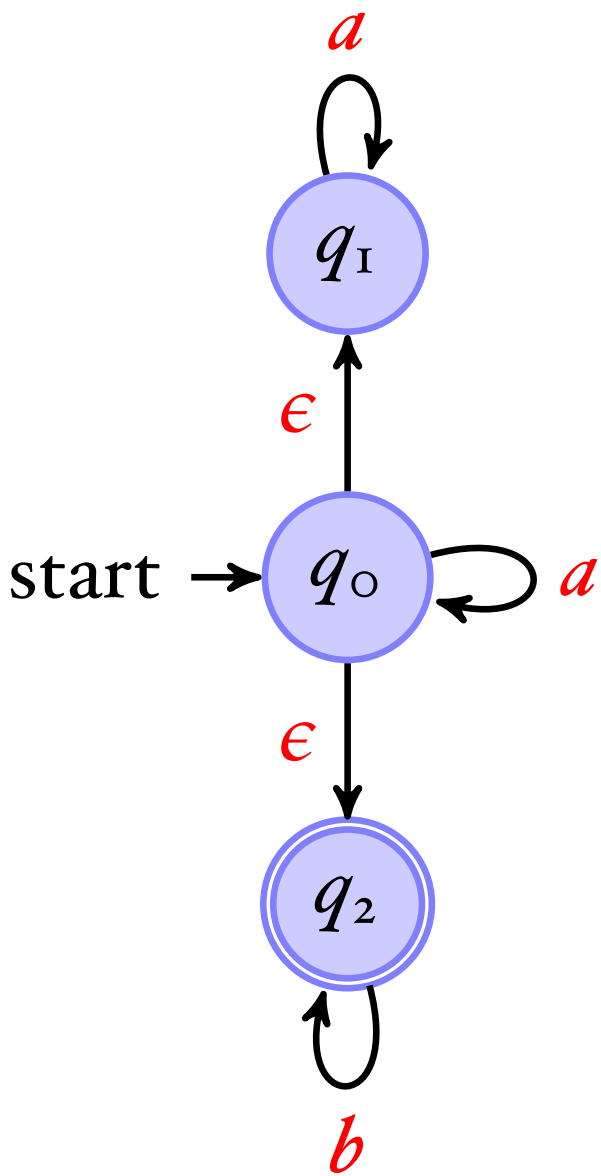

nodes	a	b
$\varnothing$	$\varnothing$	$\varnothing$
$\{0\}$	$\{0,1,2\}$	$\{2\}$
$\{1\}$	$\{1\}$	$\varnothing$
$\{2\}$	$\varnothing$	$\{2\}$
$\{0,1\}$	$\{0,1,2\}$	$\{2\}$
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$\{1,2\}$	$\{1\}$	$\{2\}$
$\{0,1,2\}$	$\{0,1,2\}$	$\{2\}$

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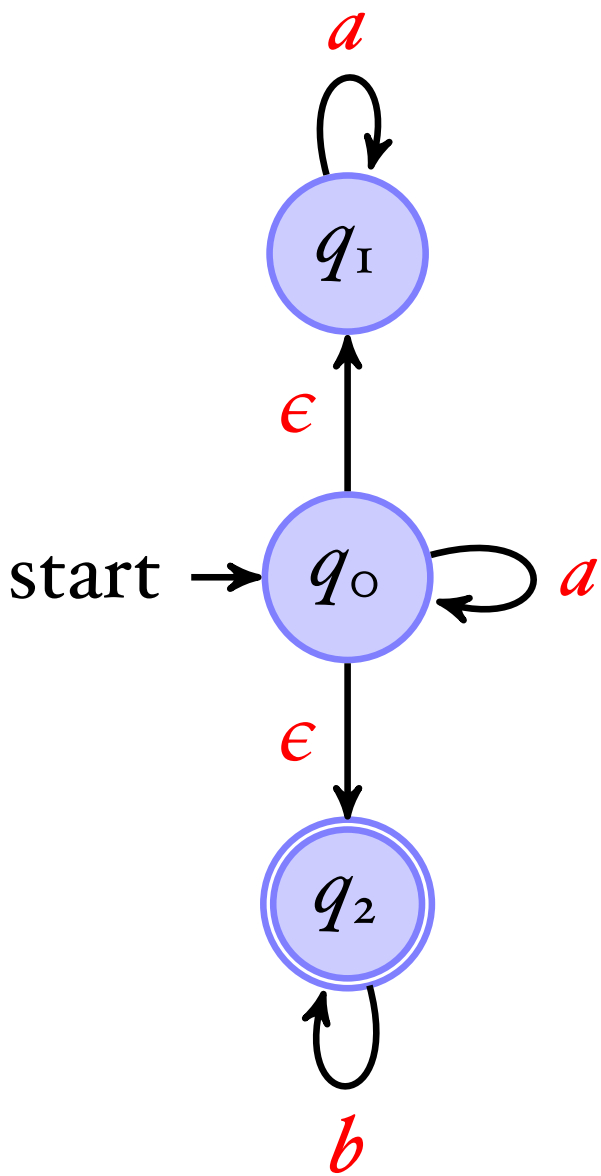


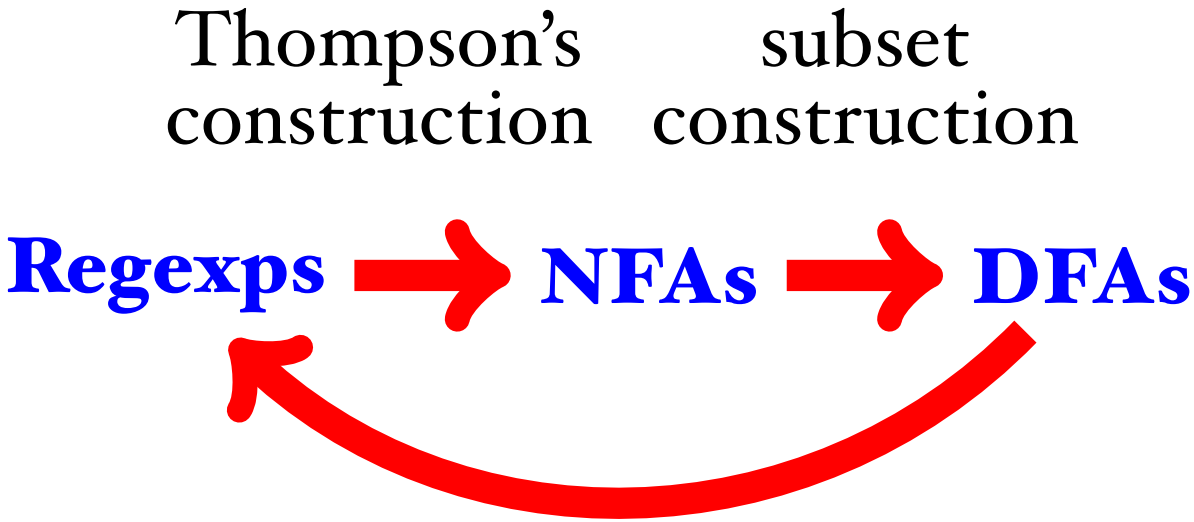
nodes		a	b
$\varnothing$		$\varnothing$	$\varnothing$
$\{0\}$		$\{0,1,2\}$	$\{2\}$
$\{1\}$		$\{1\}$	$\varnothing$
$\{2\}$	*	$\varnothing$	$\{2\}$
$\{0,1\}$		$\{0,1,2\}$	$\{2\}$
$\{0,2\}$	*	$\{0,1,2\}$	$\{2\}$
$\{1,2\}$	*	$\{1\}$	$\{2\}$
s: $\{0,1,2\}$	*	$\{0,1,2\}$	$\{2\}$

Regexps and Automata

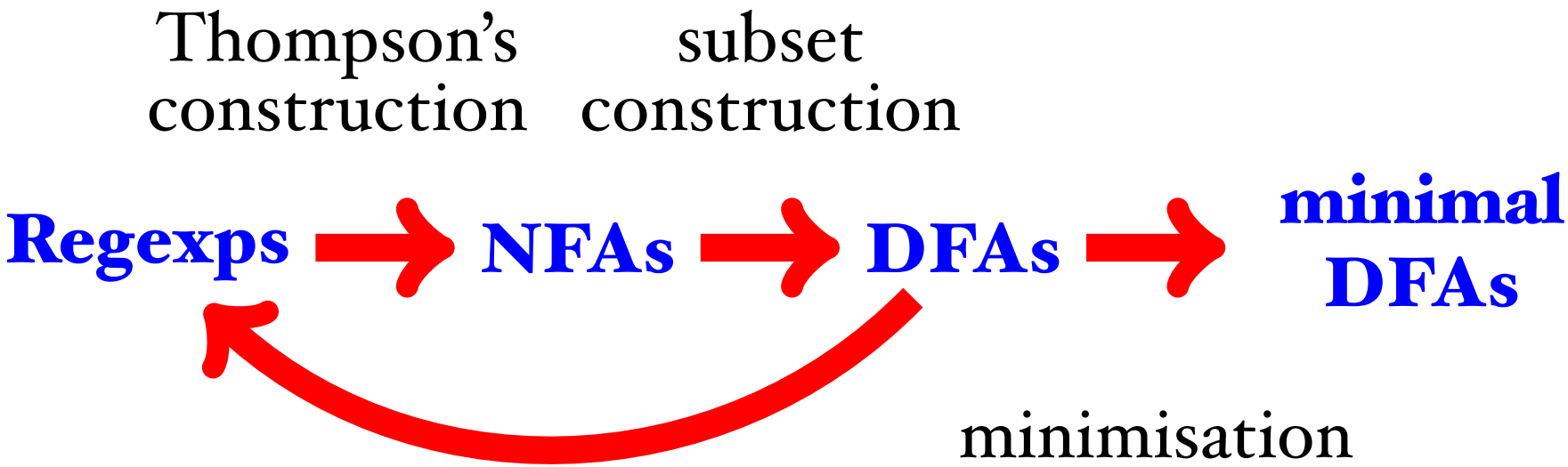
Thompson's construction

subset construction

Regexps and Automata



Regexps and Automata



Regular Languages

A language is regular iff there exists a regular expression that recognises all its strings.

or **equivalently**

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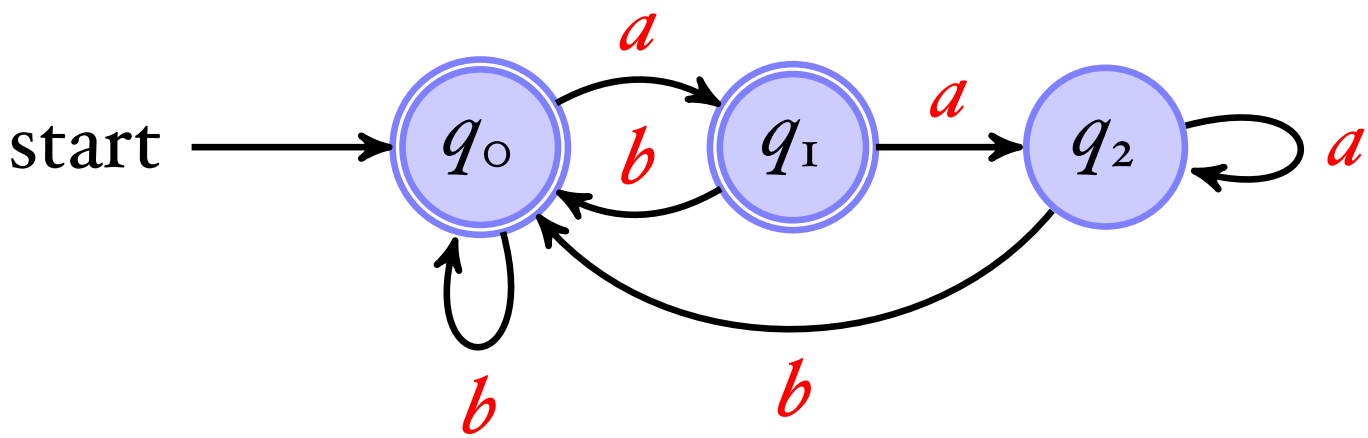
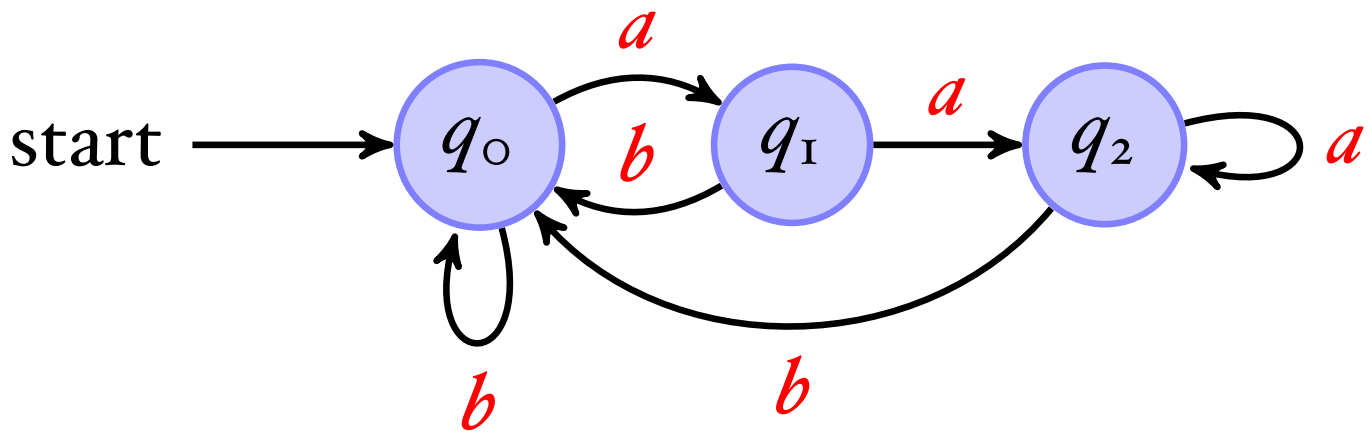
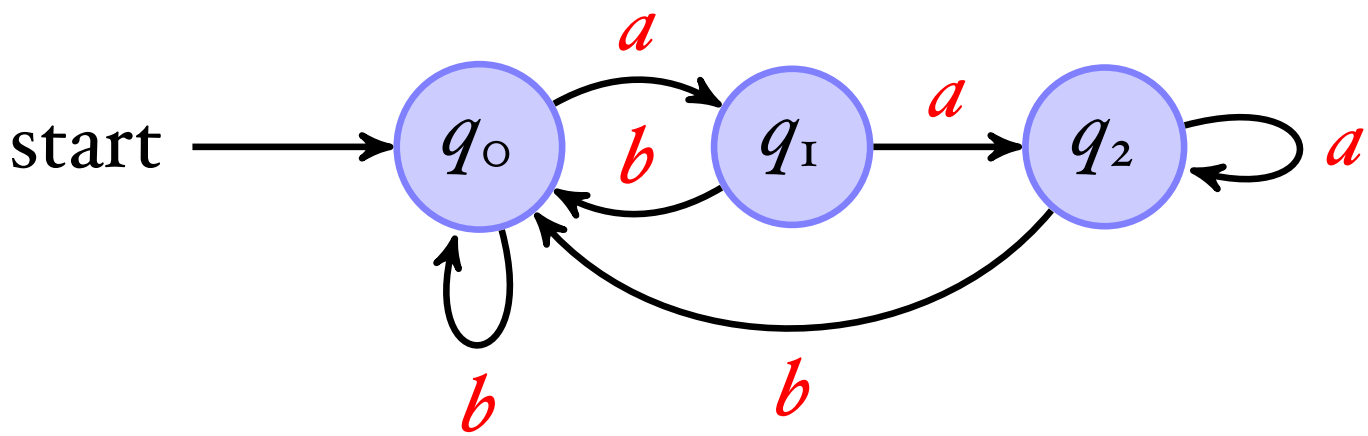
Why is every finite set of strings a regular language?

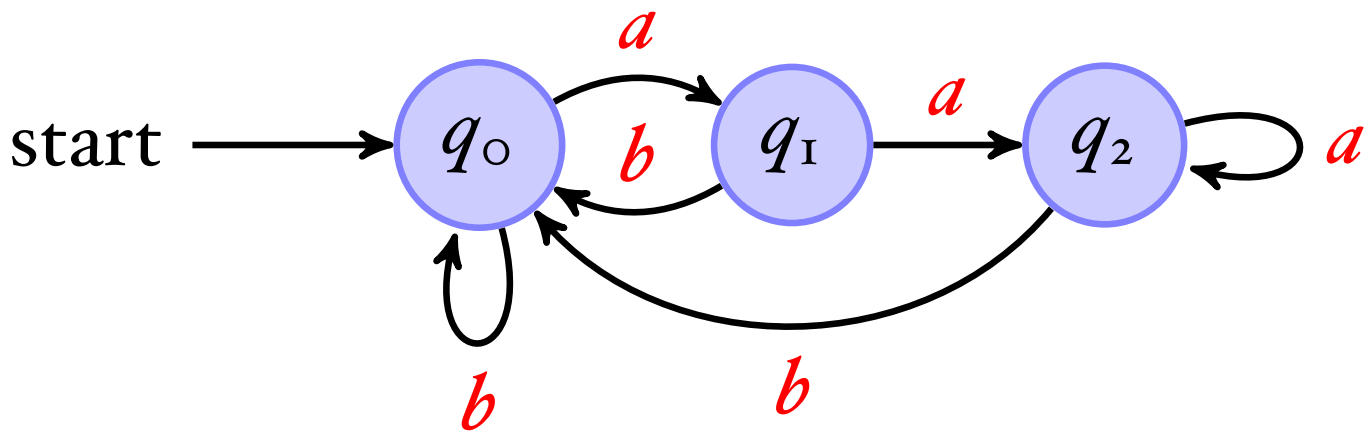
DFA to Rexp

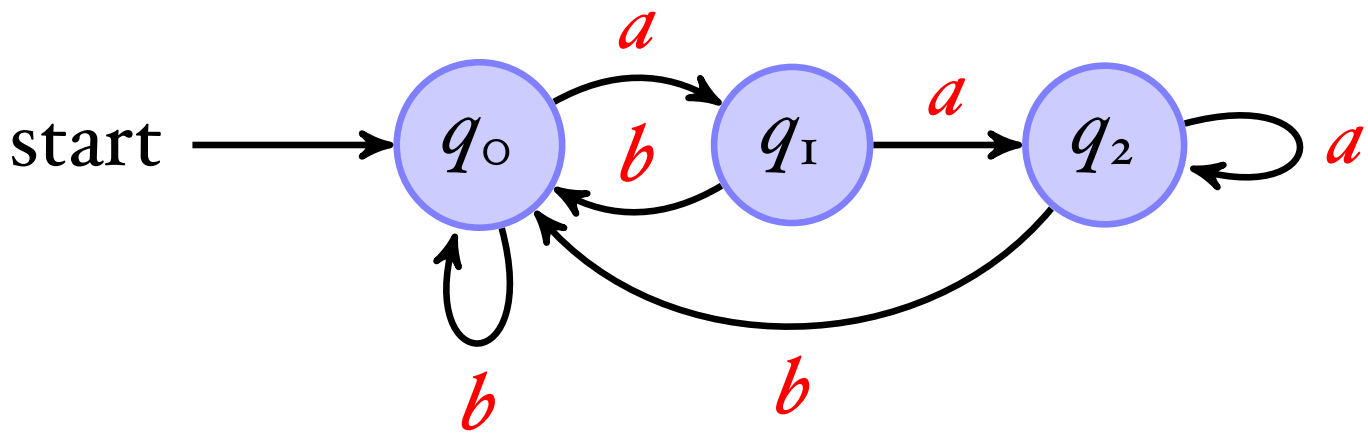



start
q_0
q_1
q_2
a
b
a
a
b
b

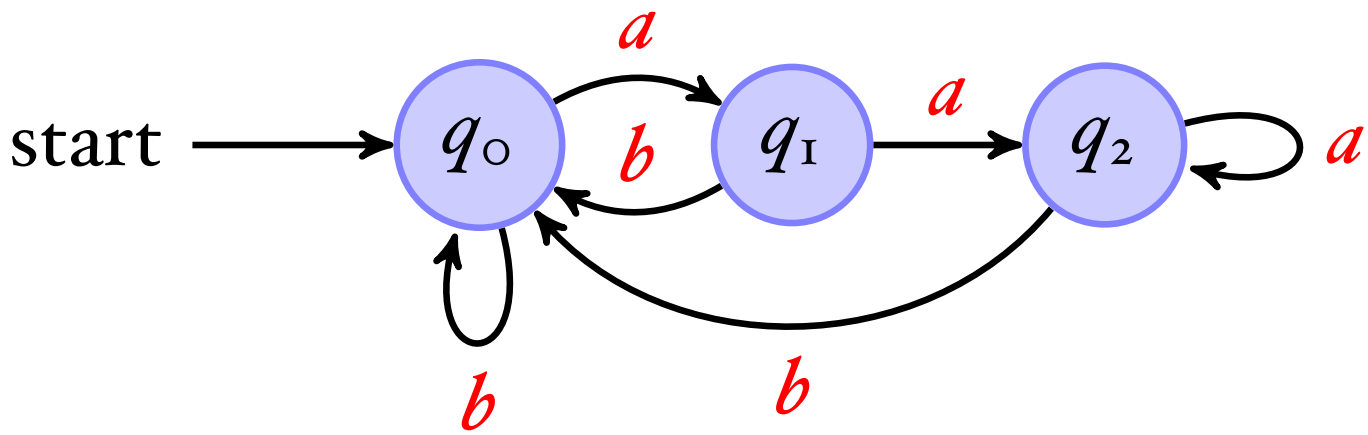


$$
\begin{aligned}
q_0 &= 2q_0 + 3q_1 + 4q_2 \\
q_1 &= 2q_0 + 3q_1 + 1q_2 \\
q_2 &= 1q_0 + 5q_1 + 2q_2
\end{aligned}
$$


start
q_0
q_1
q_2
a
a
a
b
b
b



$$q_0 = \epsilon + q_0 b + q_1 b + q_2 b$$
$$q_1 = q_0 a$$
$$q_2 = q_1 a + q_2 a$$



$$q_0 = \epsilon + q_0 b + q_1 b + q_2 b$$
$$q_1 = q_0 a$$
$$q_2 = q_1 a + q_2 a$$

Arden's Lemma:

$$\text{If } q = qr + s \text{ then } q = sr^*$$

Given the function

$$\begin{array}{rcl} rev(\varnothing) & \stackrel{\text{def}}{=} & \varnothing \\ rev(\epsilon) & \stackrel{\text{def}}{=} & \epsilon \\ rev(c) & \stackrel{\text{def}}{=} & c \\ rev(r_1 + r_2) & \stackrel{\text{def}}{=} & rev(r_1) + rev(r_2) \\ rev(r_1 \cdot r_2) & \stackrel{\text{def}}{=} & rev(r_2) \cdot rev(r_1) \\ rev(r^*) & \stackrel{\text{def}}{=} & rev(r)^* \end{array}$$

and the set

$$Rev\, A \stackrel{\text{def}}{=} \{s^{-1} \mid s \in A\}$$

prove whether

$$L(rev(r)) = Rev(L(r))$$