Homework 7

Please submit your solutions via email. Please submit only PDFs! Every solution should be preceded by the corresponding question text, like:

Qn: ...a difficult question from me...

A: ...an answer from you ... Qn + 1 ...another difficult question...

A: ...another brilliant answer from you...

Solutions will only be accepted until 20th December! Please send only one homework per email.

1. Suppose the context-sensitive grammar

$$\begin{array}{ccc}
S & \to & bSAA \mid \epsilon \\
A & \to & a \\
bA & \to & Ab
\end{array}$$

where *S* is the starting symbol of the grammar. Give a derivation of the string *"aaabaaabb"*. What can you say about the number of as and bs in the strings recognised by this grammar.

2. Consider the following grammar

$$S ::= N \cdot P$$

$$P ::= V \cdot N$$

$$N ::= N \cdot N$$

$$N ::= A \cdot N$$

$$N ::= exttt{student} \mid exttt{trainer} \mid exttt{team} \mid exttt{trains}$$

$$V ::= \mathtt{trains} \mid \mathtt{team}$$

$$A ::=$$
The $|$ the

where *S* is the start symbol and *S*, *P*, *N*, *V* and *A* are non-terminals. Using the CYK-algorithm, check whether or not the following string can be parsed by the grammar:

The trainer trains the student team

3. Transform the grammar

$$\begin{array}{ccc} A & \rightarrow & 0A1 \mid BB \\ B & \rightarrow & \epsilon \mid 2B \end{array}$$

into Chomsky normal form.

4. Consider the following grammar *G*

$$\begin{array}{l} S \rightarrow \mathtt{if0} \cdot E \cdot \mathtt{then} \cdot S \\ S \rightarrow \mathtt{print} \cdot S \\ S \rightarrow \mathtt{begin} \cdot B \cdot \mathtt{end} \\ B \rightarrow S \cdot \mathtt{;} \\ B \rightarrow S \cdot \mathtt{;} \cdot B \\ S \rightarrow num \\ E \rightarrow num \end{array}$$

 $B \rightarrow num$

where *S* is the start symbol and *S*, *E* and *B* are non-terminals.

Check each rule below and decide whether, when added to *G*, the combined grammar is ambiguous. If yes, give a string that has more than one parse tree.

- (i) $S \rightarrow \text{ifO} \cdot E \cdot \text{then} \cdot S \cdot \text{else} \cdot S$
- (ii) $B \rightarrow B \cdot B$
- (iii) $E \rightarrow (\cdot E \cdot)$
- (iv) $E \rightarrow E \cdot + \cdot E$