

# Compilers and Formal Languages (I)

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# The Goal of this Course

## Write A Compiler



# The Goal of this Course

lexer input: a string

```
"read(n);"
```

lexer output: a sequence of tokens

```
key(read); lpar; id(n); rpar; semi
```



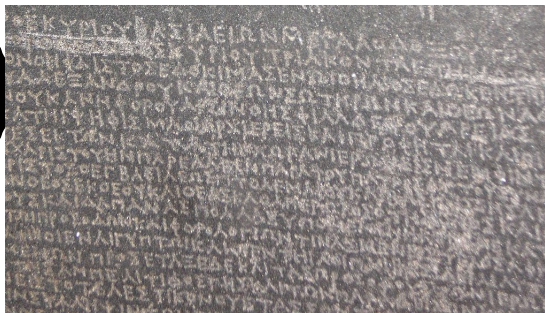
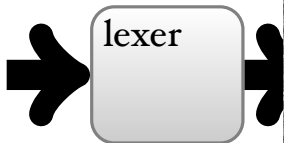
# The Goal of this Course

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"read(n);"
```

lexer output: a sequence of tokens

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key(read); lpar; id(n); rpar; semi
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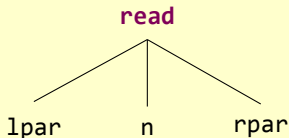


lexing  $\Rightarrow$  recognising words (Stone of Rosetta)

# The Goal of this Course

parser input: a sequence of token

parser output: an abstract syntax tree

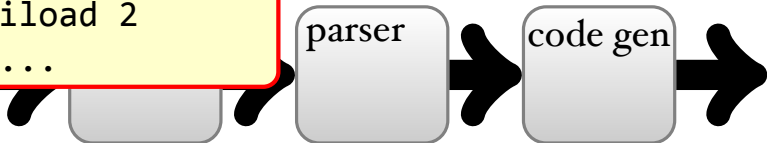


# The Goal of this Course

code generator:

```
istore 2  
iload 2  
ldc 10  
isub  
ifeq Label2  
iload 2  
...
```

## Building A Compiler



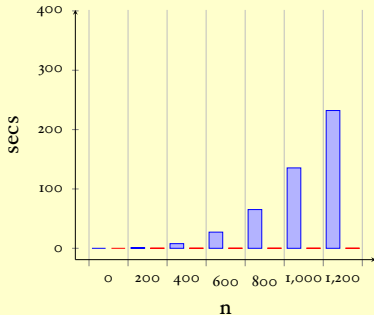
# The Goal of this Course

code generator:

```
istore 2
iload 2
ldc 10
isub
ifeq Label2
iload 2
...
```

## Building a Compiler

parse



# The subject is quite old

- Turing Machines, 1936
- Regular Expressions, 1956
- The first compiler for COBOL, 1957  
(Grace Hopper)
- But surprisingly research papers are still published nowadays



Grace Hopper

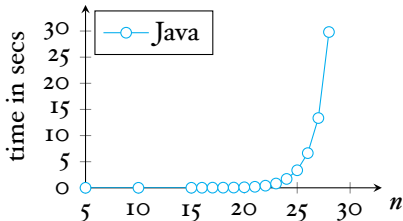
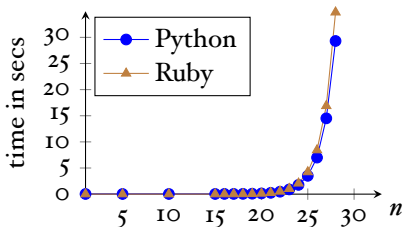
(she made it to David Letterman's Tonight Show,

<http://www.youtube.com/watch?v=aZ0xtURhfEU>)

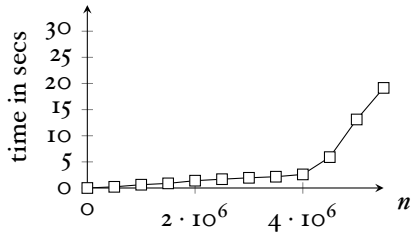
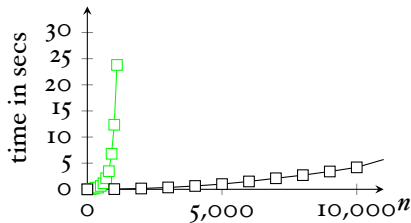


# Why Bother?

Ruby, Python, Java



Us (after next lecture)



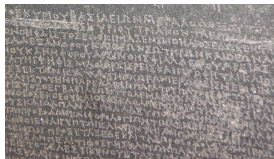
matching  $[a?]\{n\}[a]\{n\}$  and  $[a^*]^*b$  against  $\underbrace{a\dots a}_n$

# Lectures 1 - 5

transforming strings into structured data

**Lexing** based on regular expressions  
(recognising “words”)

**Parsing**  
(recognising “sentences”)



Stone of Rosetta

# Familiar Regular Expr.

`[a-zA-Z0-9_.-]+ @ [a-zA-Z0-9_.-]+ . [a-zA-Z.]{2,6}`

<code>re*</code>	matches 0 or more times
<code>re+</code>	matches 1 or more times
<code>re?</code>	matches 0 or 1 times
<code>re{n}</code>	matches exactly n number of times
<code>re{n,m}</code>	matches at least n and at most m times
<code>[...]</code>	matches any single character inside the brackets
<code>[^...]</code>	matches any single character not inside the brackets
<code>a-zA-Z</code>	character ranges
<code>\d</code>	matches digits; equivalent to <code>[0-9]</code>
<code>.</code>	matches every character except newline
<code>(re)</code>	groups regular expressions and remembers the matched text

# Today

- While the ultimate goal is to implement a small compiler (a really small one for the JVM)...

Let's start with:

- a web-crawler
- an email harvester
- (a web-scraper)

# A Web-Crawler

- 1 given an URL, read the corresponding webpage
- 2 extract all links from it
- 3 call the web-crawler again for all these links

# A Web-Crawler

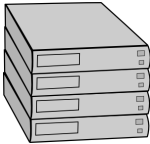
- ➊ given an URL, read the corresponding webpage
- ➋ if not possible print, out a problem
- ➌ if possible, extract all links from it
- ➍ call the web-crawler again for all these links

# A Web-Crawler

- 1 given an URL, read the corresponding webpage
- 2 if not possible print, out a problem
- 3 if possible, extract all links from it
- 4 call the web-crawler again for all these links

(we need a bound for the number of recursive calls)

(the purpose is to check all links on my own webpage)



Server

GET request



webpage



POST data



Browser



# Scala

A simple Scala function for reading webpages:

```
import io.Source
```

```
def get_page(url: String) : String = {  
  Source.fromURL(url).take(10000).mkString  
}
```

# Scala

A simple Scala function for reading webpages:

```
import io.Source
```

```
def get_page(url: String) : String = {  
  Source.fromURL(url).take(10000).mkString  
}
```

```
get_page("""http://www.inf.kcl.ac.uk/staff/urbanc/""")
```

# Scala

A simple Scala function for reading webpages:

```
import io.Source
```

```
def get_page(url: String) : String = {  
  Source.fromURL(url).take(10000).mkString  
}
```

```
get_page("""http://www.inf.kcl.ac.uk/staff/urbanc/""")
```

A slightly more complicated version for handling errors:

```
def get_page(url: String) : String = {  
  Try(Source.fromURL(url).take(10000).mkString).  
  getOrElse { println(s" Problem with: $url"); ""}  
}
```

# A Regular Expression

- ... is a pattern or template for specifying strings

```
"https?://[^\"]*"
```

matches for example

```
"http://www.foobar.com"
```

```
"https://www.tls.org"
```

# A Regular Expression

- ... is a pattern or template for specifying strings

```
""""https?://[^\"]*"""".r
```

matches for example

```
”http://www.foobar.com”
```

```
”https://www.tls.org”
```

# Finding Operations

**rexp.findAllIn(string)**

returns a list of all (sub)strings that match the regular expression

**rexp.findFirstIn(string)**

returns either

- None if no (sub)string matches or
- Some(s) with the first (sub)string

```
val http_pattern = """https?://[^\"]*""".r

def unquote(s: String) = s.drop(1).dropRight(1)

def get_all_URLs(page: String) : Set[String] =
  http_pattern.findAllIn(page).map(unquote).toSet

def crawl(url: String, n: Int) : Unit = {
  if (n == 0) ()
  else {
    println(s"Visiting: $n $url")
    for (u <- get_all_URLs(get_page(url))) crawl(u, n - 1)
  }
}

crawl(some_start_URL, 2)
```

A version that only crawls links in “my” domain:

```
val my_urls = """urbanc""".r

def crawl(url: String, n: Int) : Unit = {
  if (n == 0) ()
  else if (my_urls.findFirstIn(url) == None) {
    println(s"Visiting: $n $url")
    get_page(url); ()
  }
  else {
    println(s"Visiting: $n $url")
    for (u <- get_all_URLs(get_page(url))) crawl(u, n - 1)
  }
}
```



## A little email harvester:

```
val http_pattern = """https?://[^\"]*""".r
val email_pattern =
  """([a-z0-9_\. -]+)@([\da-z\.-]+)\.([a-z\.]{2,6})""".r

def print_str(s: String) =
  if (s == "") () else println(s)

def crawl(url: String, n: Int) : Unit = {
  if (n == 0) ()
  else {
    println(s"Visiting: $n $url")
    val page = get_page(url)
    print_str(email_pattern.findAllIn(page).mkString("\n"))
    for (u <- get_all_URLs(page).par) crawl(u, n - 1)
  }
}
```

<http://net.tutsplus.com/tutorials/other/8-regular-expressions-you-should-know/>

# Regular Expressions

Their inductive definition:

$r ::=$	<b>0</b>	nothing
	<b>1</b>	empty string / "" / []
	$c$	character
	$r_1 + r_2$	alternative / choice
	$r_1 \cdot r_2$	sequence
	$r^*$	star (zero or more)

Th

```
abstract class Rexp
case object ZERO extends Rexp
case object ONE extends Rexp
case class CHAR(c: Char) extends Rexp
case class ALT(r1: Rexp, r2: Rexp) extends Rexp
case class SEQ(r1: Rexp, r2: Rexp) extends Rexp
case class STAR(r: Rexp) extends Rexp
```

$r ::= \mathbf{0}$	nothing
$\mathbf{I}$	empty string / "" / []
$c$	character
$r_1 + r_2$	alternative / choice
$r_1 \cdot r_2$	sequence
$r^*$	star (zero or more)

# Strings

...are lists of characters. For example "hello"

$[h, e, l, l, o]$  or just *hello*

the empty string:  $[]$  or ""

the concatenation of two strings:

$s_1 @ s_2$

$foo @ bar = foobar, baz @ [] = baz$

# Languages, Strings

- **Strings** are lists of characters, for example

$[], abc$  (Pattern match:  $c::s$ )

- A **language** is a set of strings, for example

$\{[], bello, foobar, a, abc\}$

- **Concatenation** of strings and languages

$foo @ bar = foobar$

$A @ B \stackrel{\text{def}}{=} \{s_1 @ s_2 \mid s_1 \in A \wedge s_2 \in B\}$

# The Meaning of a Regular Expression

$$L(\mathbf{0}) \stackrel{\text{def}}{=} \{\}$$

$$L(\mathbf{I}) \stackrel{\text{def}}{=} \{\emptyset\}$$

$$L(c) \stackrel{\text{def}}{=} \{[c]\}$$

$$L(r_1 + r_2) \stackrel{\text{def}}{=} L(r_1) \cup L(r_2)$$

$$L(r_1 \cdot r_2) \stackrel{\text{def}}{=} \{s_1 @ s_2 \mid s_1 \in L(r_1) \wedge s_2 \in L(r_2)\}$$

$$L(r^*) \stackrel{\text{def}}{=}$$

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$$L(r)^{n+1} \stackrel{\text{def}}{=} L(r) @ L(r)^n$$

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$$\{s_1 @ s_2 \mid s_1 \in L(r) \wedge s_2 \in L(r)^n\}$$



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$$L(r^*) \stackrel{\text{def}}{=} \bigcup_{0 \leq n} L(r)^n$$

$$L(r)^\circ \stackrel{\text{def}}{=} \{\emptyset\}$$

$$L(r)^{n+1} \stackrel{\text{def}}{=} L(r) @ L(r)^n \quad (\text{append on sets})$$
$$\{s_1 @ s_2 \mid s_1 \in L(r) \wedge s_2 \in L(r)^n\}$$

# The Meaning of Matching

A regular expression  $r$  matches a string  $s$  provided

$$s \in L(r)$$

...and the point of the next lecture is to decide this problem as fast as possible (unlike Python, Ruby, Java)

# Written Exam

- Accounts for 80%.
- You will understand the question “*Is this relevant for the exam?*” is very demotivating for the lecturer!
- Deal: Whatever is in the homework (and is not marked “*optional*”) is relevant for the exam.
- Each lecture has also a handout. There are also handouts about notation and Scala.

# Coursework

- Accounts for 20%. Two strands. Choose **one**!

## Strand 1

- four programming tasks:
  - matcher (4%, 19.10.)
  - lexer (5%, 03.11.)
  - parser (5%, 23.11.)
  - compiler (6%, 7.12.)

## Strand 2

- one task: prove the correctness of a regular expression matcher in the Isabelle theorem prover
- 20%, submission 7.12.

- Solving more than one strand will **not** give you more marks.

# Lecture Capture

- Hope it works...

# Lecture Capture

- Hope it works...
- It is important to use lecture capture wisely:
  - Lecture recordings are a study and revision aid.
  - Statistically, there is a clear and direct link between attendance and attainment: Students who do not attend lectures, do less well in exams.
- Attending a lecture is more than watching it online – if you do not attend, you miss out!

**Questions?**