

Homework 5

1. Consider the basic regular expressions

$$r ::= \emptyset \mid \epsilon \mid c \mid r_1 + r_2 \mid r_1 \cdot r_2 \mid r^*$$

and suppose you want to show a property $P(r)$ for all regular expressions r by structural induction. Write down which cases do you need to analyse. State clearly the induction hypotheses if applicable in a case.

2. Define a regular expression, written *ALL*, that can match every string. This definition should be in terms of the following extended regular expressions:

$$r ::= \emptyset \mid \epsilon \mid c \mid r_1 + r_2 \mid r_1 \cdot r_2 \mid r^* \mid \sim r$$

3. Assume the delimiters for comments are */** and **/*. Give a regular expression that can recognise comments of the form

$$/* \dots */$$

where the three dots stand for arbitrary characters, but not comment delimiters.

4. Define the following regular expressions

r^+	(one or more matches)
$r^?$	(zero or one match)
$r^{\{n\}}$	(exactly n matches)
$r^{\{m,n\}}$	(at least m and maximal n matches, with the assumption $m \leq n$)

in terms of the usual basic regular expressions

$$r ::= \emptyset \mid \epsilon \mid c \mid r_1 + r_2 \mid r_1 \cdot r_2 \mid r^*$$

5. Give the regular expressions for lexing a language consisting of identifiers, left-parenthesis *(*, right-parenthesis *)*, numbers that can be either positive or negative, and the operations *+*, *-* and ***.

Decide whether the following strings can be lexed in this language?

- (a) "(a3+3)*b"
- (b) ")()++-33"
- (c) "(b42/3)*3"

In case they can, give the corresponding token sequences. (Hint: Observe the maximal munch rule and the priorities of your regular expressions that make the process of lexing unambiguous.)

6. (Optional) Recall the definitions for Der and der from the lectures. Prove by induction on r the property that

$$L(der\ c\ r) = Der\ c\ (L(r))$$

holds.