Homework 2

Please submit your solutions via email. Please submit only ASCII text or PDFs. Every solution should be preceded by the corresponding question, like:

> **Q***n*:   ...a difficult question from me...
> **A:**   ...an answer from you ...
> **Q***n* + 1   ...another difficult question...
> **A:**   ...another brilliant answer from you...

Solutions will only be accepted until 30th December! Please send only one homework per email.

1. What is the language recognised by the regular expressions $(\mathbf{0}^*)^*$.

2. Review the first handout about sets of strings and read the second handout. Assuming the alphabet is the set $\{a, b\}$, decide which of the following equations are true in general for arbitrary languages A, B and C:

$$\begin{aligned}
(A \cup B)@C &\overset{?}{=} A@C \cup B@C \\
A^* \cup B^* &\overset{?}{=} (A \cup B)^* \\
A^*@A^* &\overset{?}{=} A^* \\
(A \cap B)@C &\overset{?}{=} (A@C) \cap (B@C)
\end{aligned}$$

   In case an equation is true, give an explanation; otherwise give a counter-example.

3. Given the regular expressions $r_1 = \mathbf{1}$ and $r_2 = \mathbf{0}$ and $r_3 = a$. How many strings can the regular expressions r_1^*, r_2^* and r_3^* each match?

4. Give regular expressions for (a) decimal numbers and for (b) binary numbers. (Hint: Observe that the empty string is not a number. Also observe that leading 0s are normally not written.)

5. Decide whether the following two regular expressions are equivalent $(\mathbf{1} + a)^* \overset{?}{\equiv} a^*$ and $(a \cdot b)^* \cdot a \overset{?}{\equiv} a \cdot (b \cdot a)^*$.

6. Given the regular expression $r = (a \cdot b + b)^*$. Compute what the derivative of r is with respect to a, b and c. Is r nullable?

7. Prove that for all regular expressions r we have

$$\textit{nullable}(r) \quad \text{if and only if} \quad [] \in L(r)$$

Write down clearly in each case what you need to prove and what are the assumptions.

8. Define what is meant by the derivative of a regular expressions with respect to a character. (Hint: The derivative is defined recursively.)

9. Assume the set *Der* is defined as

$$Der\ c\ A \stackrel{\text{def}}{=} \{s \mid c{::}s \in A\}$$

What is the relation between *Der* and the notion of derivative of regular expressions?

10. Give a regular expression over the alphabet $\{a, b\}$ recognising all strings that do not contain any substring bb and end in a.

11. Do $(a + b)^* \cdot b^+$ and $(a^* \cdot b^+) + (b^* \cdot b^+)$ define the same language?

12. Define the function *zeroable* by recursion over regular expressions. This function should satisfy the property

$$zeroable(r) \text{ if and only if } L(r) = \{\} \qquad (*)$$

The function *nullable* for the not-regular expressions can be defined by

$$nullable(\sim r) \stackrel{\text{def}}{=} \neg(nullable(r))$$

Unfortunately, a similar definition for *zeroable* does not satisfy the property in $(*)$:

$$zeroable(\sim r) \stackrel{\text{def}}{=} \neg(zeroable(r))$$

Find out why?

13. Give a regular expressions that can recognise all strings from the language $\{a^n \mid \exists k.\ n = 3k + 1\}$.