

Homework 2

Please submit your solutions via email. Please submit only ASCII text or PDFs. Every solution should be preceded by the corresponding question, like:

Q_n: ...a difficult question from me...
A: ...an answer from you ...
Q_n + 1 ...another difficult question...
A: ...another brilliant answer from you...

Solutions will only be accepted until 30th December! Please send only one homework per email.

1. What is the language recognised by the regular expressions $(0^*)^*$.
2. Review the first handout about sets of strings and read the second handout. Assuming the alphabet is the set $\{a, b\}$, decide which of the following equations are true in general for arbitrary languages A, B and C :

$$\begin{aligned}(A \cup B)@C & \stackrel{?}{=} A@C \cup B@C \\ A^* \cup B^* & \stackrel{?}{=} (A \cup B)^* \\ A^*@A^* & \stackrel{?}{=} A^* \\ (A \cap B)@C & \stackrel{?}{=} (A@C) \cap (B@C)\end{aligned}$$

In case an equation is true, give an explanation; otherwise give a counterexample.

3. Given the regular expressions $r_1 = \mathbf{1}$ and $r_2 = \mathbf{0}$ and $r_3 = a$. How many strings can the regular expressions r_1^* , r_2^* and r_3^* each match?
4. Give regular expressions for (a) decimal numbers and for (b) binary numbers. (Hint: Observe that the empty string is not a number. Also observe that leading 0s are normally not written.)
5. Decide whether the following two regular expressions are equivalent $(\mathbf{1} + a)^* \stackrel{?}{=} a^*$ and $(a \cdot b)^* \cdot a \stackrel{?}{=} a \cdot (b \cdot a)^*$.
6. Given the regular expression $r = (a \cdot b + b)^*$. Compute what the derivative of r is with respect to a, b and c . Is r nullable?
7. Prove that for all regular expressions r we have

$$\text{nullable}(r) \quad \text{if and only if} \quad \epsilon \in L(r)$$

Write down clearly in each case what you need to prove and what are the assumptions.

- Define what is meant by the derivative of a regular expressions with respect to a character. (Hint: The derivative is defined recursively.)
- Assume the set Der is defined as

$$Der\ c\ A \stackrel{\text{def}}{=} \{s \mid c::s \in A\}$$

What is the relation between Der and the notion of derivative of regular expressions?

- Give a regular expression over the alphabet $\{a, b\}$ recognising all strings that do not contain any substring bb and end in a .
- Do $(a + b)^* \cdot b^+$ and $(a^* \cdot b^+) + (b^* \cdot b^+)$ define the same language?
- Define the function *zeroable* by recursion over regular expressions. This function should satisfy the property

$$\text{zeroable}(r) \text{ if and only if } L(r) = \{\} \quad (*)$$

The function *nullable* for the not-regular expressions can be defined by

$$\text{nullable}(\sim r) \stackrel{\text{def}}{=} \neg(\text{nullable}(r))$$

Unfortunately, a similar definition for *zeroable* does not satisfy the property in (*):

$$\text{zeroable}(\sim r) \stackrel{\text{def}}{=} \neg(\text{zeroable}(r))$$

Find out why?

- Give a regular expressions that can recognise all strings from the language $\{a^n \mid \exists k. n = 3k + 1\}$.
- Give a regular expression that can recognise an odd number of *as* or an even number of *bs*.