



CSCI 742 - Compiler Construction

Lecture 26
Introduction to Code Generation
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Compiler Phases

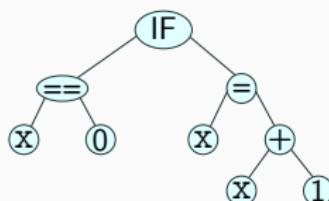
Source Code
(concrete syntax)

if (x==0) x=x+1;

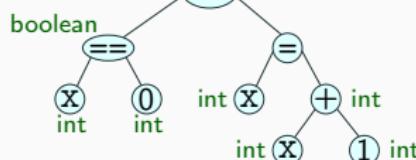
Token Stream

if (x == 0) x = x + 1 ;

Abstract Syntax Tree
(AST)

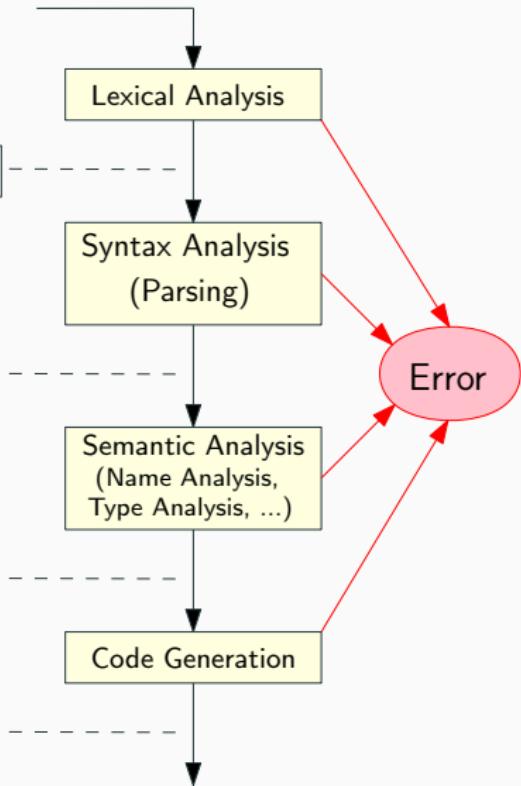


Attributed AST



Machine Code

```
16: iload_2
17: ifne 24
20: iload_2
21: iconst_1
22: iadd
23: istore_2
24: ...
```



Code Generation Example

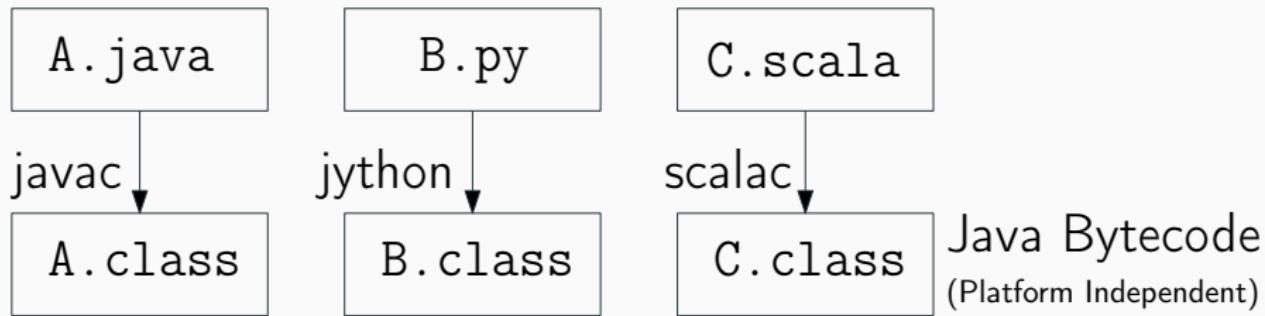
- Phase after type checking emits such bytecode instructions

```
while (i > 0) {      javac Test.java
    i += 2 * j + 1;  javap -c Test
    j = j - 5;
    System.out.println(j);
}

5: iload_1
6: ifle           31
9: iload_1
10:  iconst_2
11: iload_2
12: imul
13:  iconst_1
14: iadd
15: iadd
16: istore_1
17: iload_2
18:  iconst_5
19: isub
20: istore_2
21: getstatic      #2 // System.out
24: iload_2
25: invokevirtual #3 // println
28: goto          5
31: // ...
```

Java Virtual Machine (JVM)

- Programs are written in Java or other languages
- Compiler translates them to Java Bytecode
- Platform-specific Java Virtual Machine executes Java Bytecode

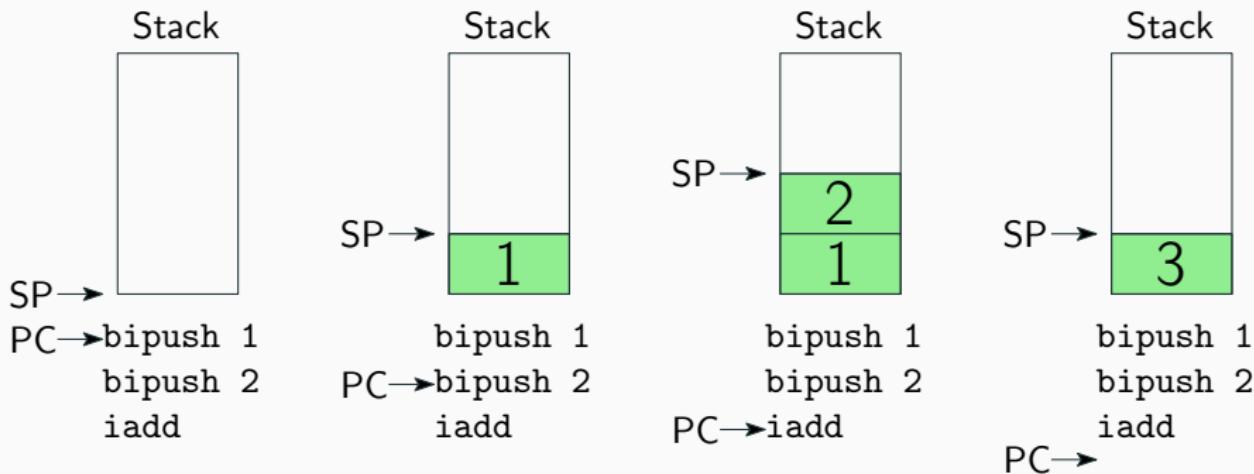


Platform-Specific Java Virtual Machine (JVM)



Java Virtual Machine (JVM)

- JVM is a stack machine: evaluation of expressions uses a stack (operand stack)
- Instructions fetch their arguments from the top of the operand stack
- Instructions store their results at the top of the operand stack



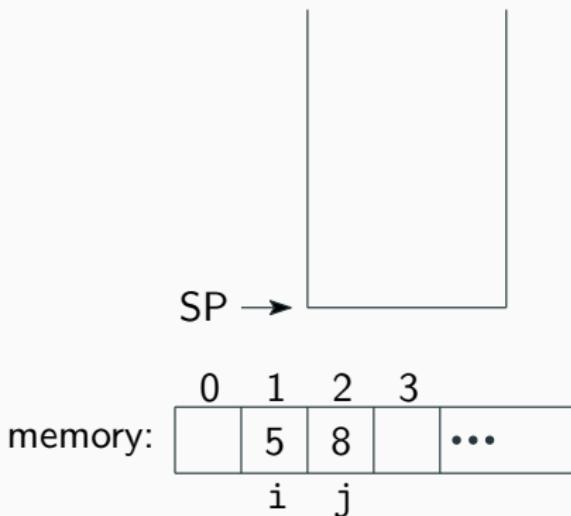
Why a Stack Machine

- A simple evaluation model: no variables or registers
- Each operation:
 - takes operands from top of stack
 - puts results back at top of stack
- Instruction “add” as opposed to “add r1, r2”
- Simpler compiler, more compact programs

Local Variables

- In memory space of a function there is an array V to store local variables and arguments
- `iload`: push the value of a local variable onto the stack
- `istore`: pop the value from the stack and store it in a local variable
- Initially the arguments x_1, \dots, x_n are stored in local variables array $V[1], \dots, V[n]$
- $V[0]$ holds the reference to this
 - object on which the method is invoked

Stack Machine Execution Example



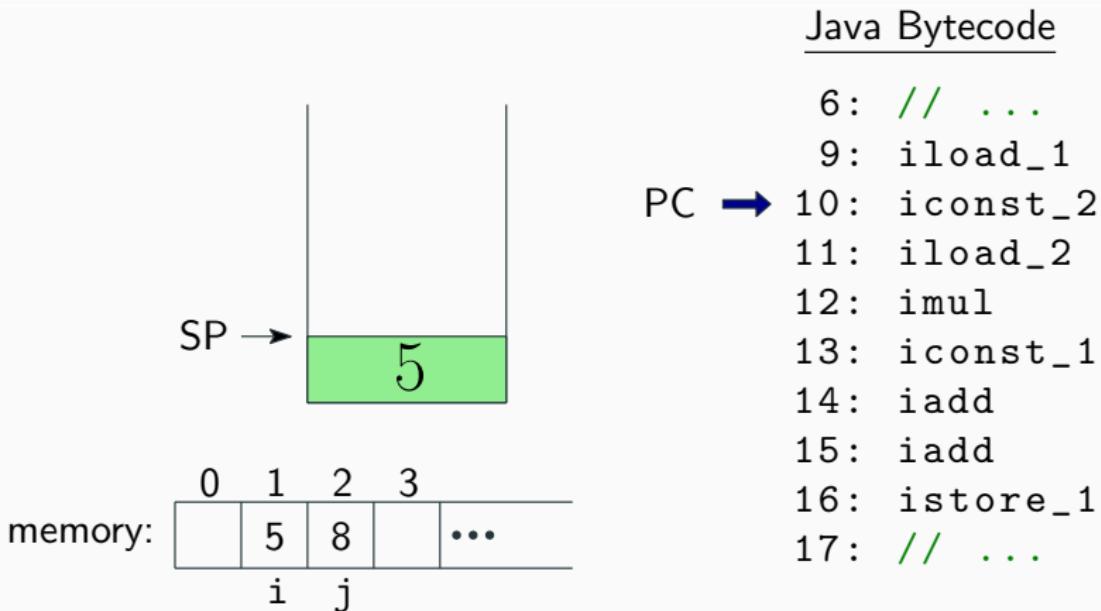
Java Bytecode

PC →

```
6: // ...
9: iload_1
10: iconst_2
11: iload_2
12: imul
13: iconst_1
14: iadd
15: iadd
16: istore_1
17: // ...
```

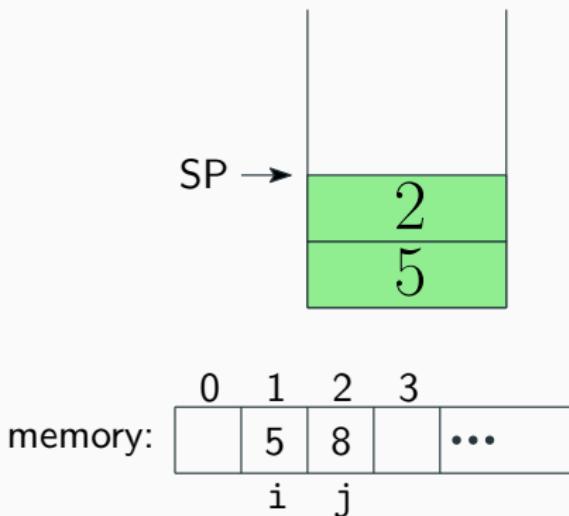
Java Statement: $i += 2 * j + 1;$

Stack Machine Execution Example



Java Statement: `i += 2 * j + 1;`

Stack Machine Execution Example

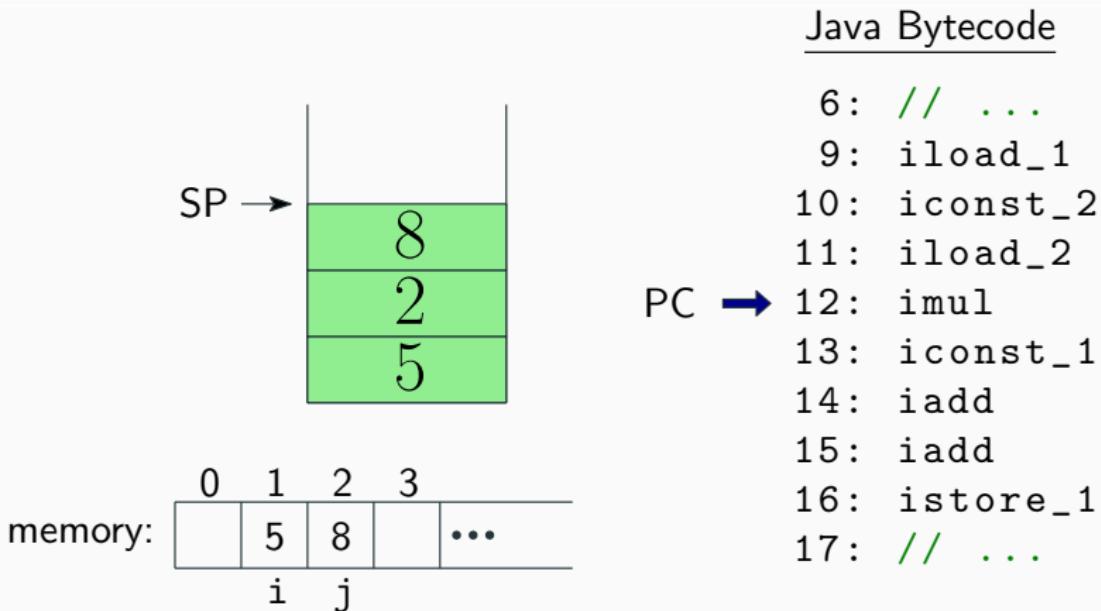


Java Bytecode

6: // ...
9: iload_1
10: iconst_2
PC → 11: iload_2
12: imul
13: iconst_1
14: iadd
15: iadd
16: istore_1
17: // ...

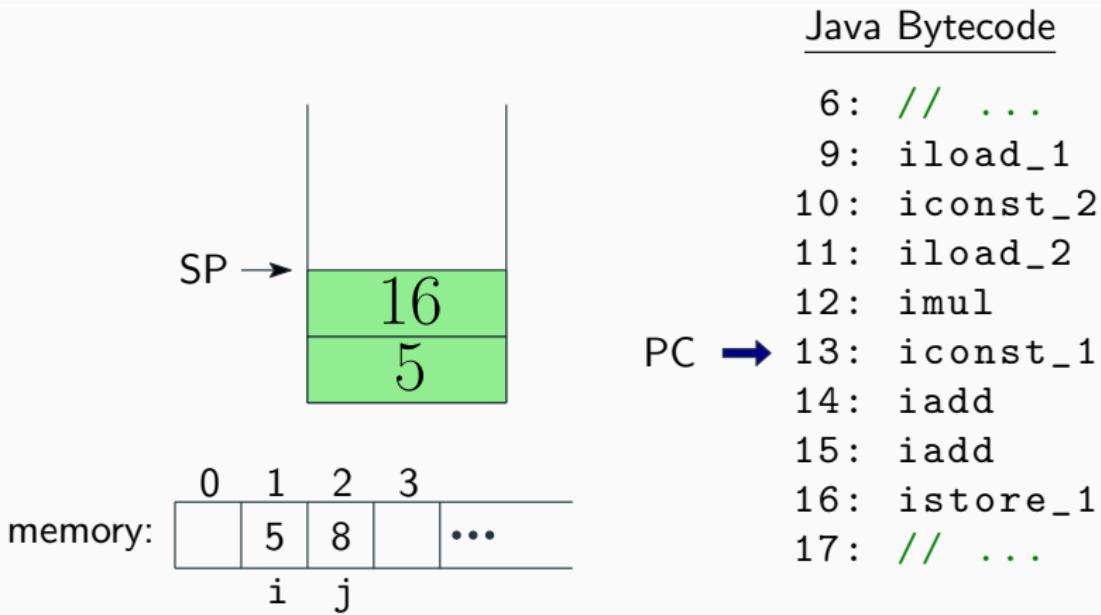
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Stack Machine Execution Example



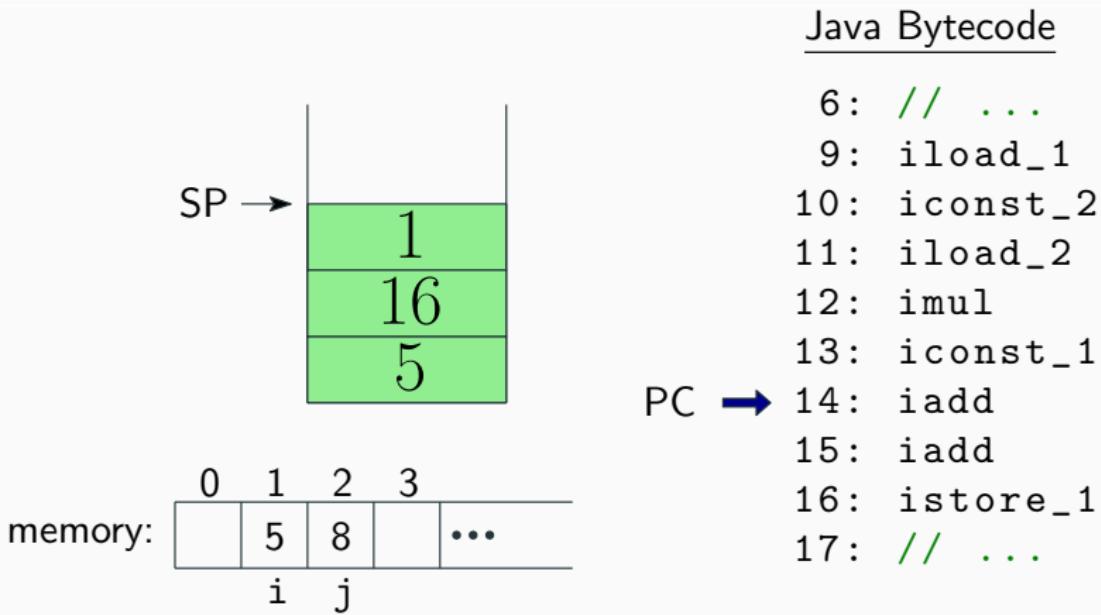
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Stack Machine Execution Example



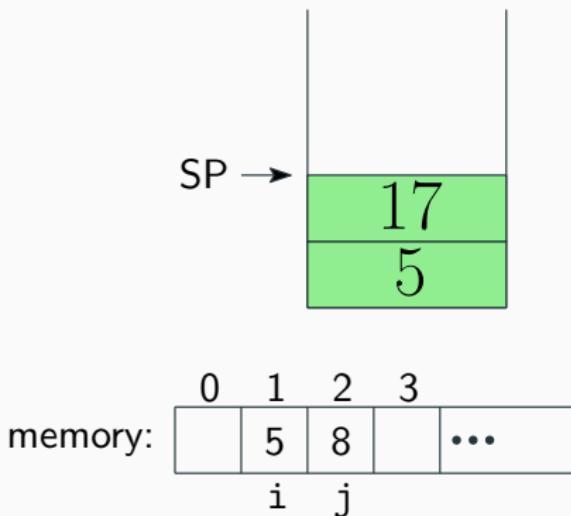
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Stack Machine Execution Example



Java Statement: `i += 2 * j + 1;`

Stack Machine Execution Example

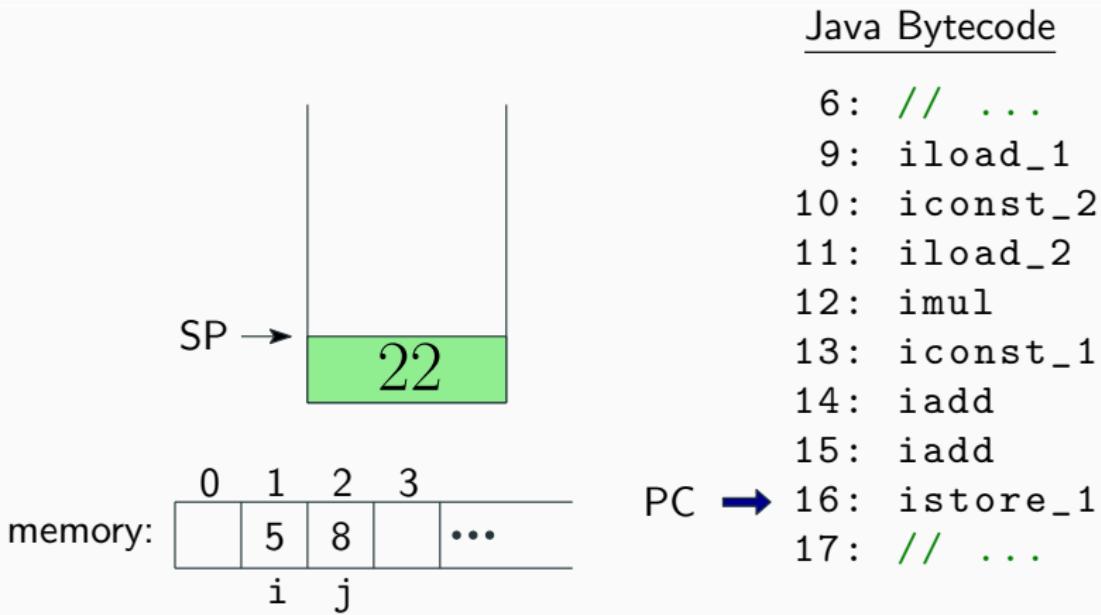


Java Bytecode

6: // ...
9: iload_1
10: iconst_2
11: iload_2
12: imul
13: iconst_1
14: iadd
PC → 15: iadd
16: istore_1
17: // ...

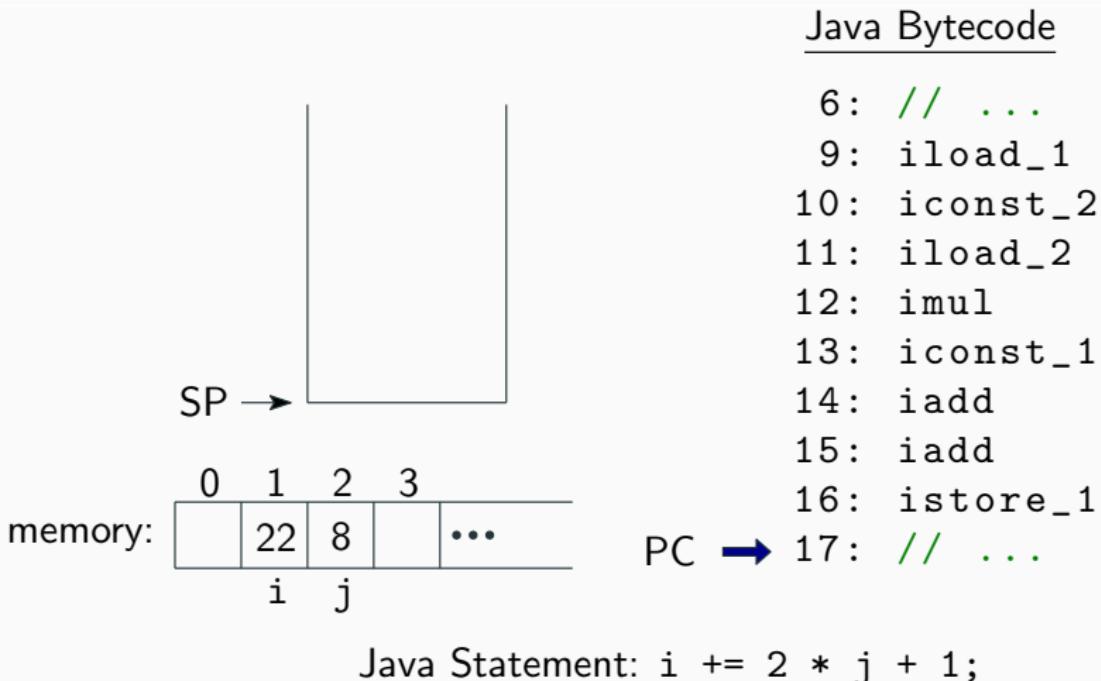
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Stack Machine Execution Example



Java Statement: `i += 2 * j + 1;`

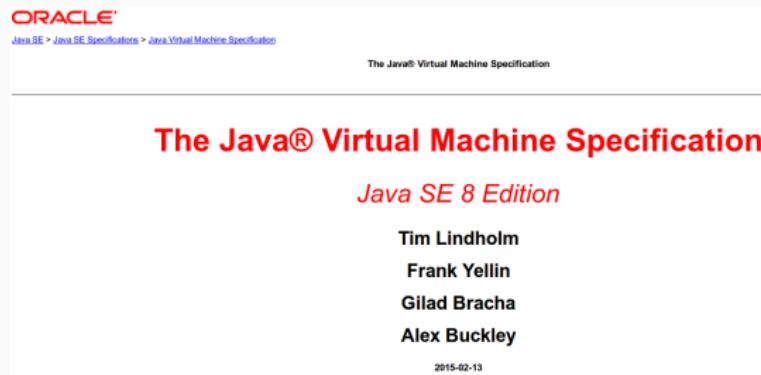
Stack Machine Execution Example



Instructions in JVM

- Separate for each type, including
 - integer types (iadd, imul, iload, istore, bipush)
 - reference types (aload, astore)
- Why are they separate if not in e.g. x86?
 - Memory safety
 - Each reference points to a valid allocated object
- Conditionals and jumps
- Further high-level operations
 - array operations
 - object method and field access

Java Virtual Machine Specifications



<http://docs.oracle.com/javase/specs/jvms/se8/html/index.html>

- Use `javac -g *.java` to compile
- Use `javap -c -l ClassName` to explore

Selected Instructions

iload_x	Loads the integer value of the local variable in slot x on the stack. $x \in \{0, 1, 2, 3\}$
iload_X	Loads the value of the local variable pointed to by index X on the top of the stack.
iconst_x	Loads the integer constant x on the stack. $x \in \{0, 1, 2, 3, 4, 5\}$
bipush_X	Like iconst, but for arbitrarily large X
istore_x	Stores the current value on top of the stack in the local variable in slot $x \in \{0, 1, 2, 3\}$
istore_X	Stores the current value on top of the stack in the local variable indexed by X.
ireturn	Method return statement (note that the return value has to have been put on the top of the stack beforehand).
iadd	Pop two (integer) values from the stack, add them and put the result back on the stack.
isub	Pop two (integer) values from the stack, subtract them and put the result back on the stack.

Selected Instructions

imult	Pop two (integer) values from the stack, multiply them and put the result back on the stack.
idiv	Pop two (integer) values from the stack, divide them and put the result back on the stack.
irem	Pop two (integer) values from the stack, put the result of $x_1 \% x_2$ back on the stack.
ineg	Negate the value on the stack.
iinc x, y	Increment the variable in slot x by amount y.
ior	Bitwise OR for the two integer values on the stack.
iand	Bitwise AND for the two integer values on the stack.
ixor	Bitwise XOR for the two integer values on the stack.
ifXX L	Pop one value from the stack, compare it zero according to the operator XX. If the condition is satisfied, jump to the instruction given by label L. XX ∈ {eq, lt, le, ne, gt, ge, null, nonnull}

Selected Instructions

if_icmpXX L	Pop two values from the stack and compare against each other. Rest as ifXX L.
goto L	Unconditional jump to instruction given by the label L.
pop	Discard word currently on top of the stack.
dup	Duplicate word currently on top of the stack.
swap	Swaps the two top values on the stack.
aload_x	Loads an object reference from slot x.
aload X	Loads an object reference from local variable indexed by X.
iaload	Loads onto the stack an integer from an array. The stack must contain the array reference and the index.
iastore	Stores an integer in an array. The stack must contain the array reference, the index and the value, in that order.

Example: Twice

```
class Expr1 {  
    public static int twice(int x) {  
        return x*2;  
    }  
}
```

```
> javac -g Expr1.java; javap -c -l Expr1
```

```
public static int twice(int);
```

Code:

```
0: iload_0 // load int from var 0 to top of stack  
1: iconst_2 // push 2 on top of stack  
2: imul     // replace two topmost elements with their product  
3: ireturn   // return top of stack
```

Example: Area

```
class Expr2 {  
    public static int cubeArea(int a,int b,int c) {  
        return (a*b + b*c + a*c) * 2;  
    } }  
  
> javac -g Expr2.java; javap -c -l Expr2
```

LocalVariableTable:

Start	Length	Slot	Name	Signature
0	14	0	a	I
0	14	1	b	I
0	14	2	c	I

```
public static int  
cubeArea(int,int,int);  
Code:  
0: iload_0  
1: iload_1  
2: imul  
3: iload_1  
4: iload_2  
5: imul  
6: iadd  
7: iload_0  
8: iload_2  
9: imul  
10: iadd  
11: iconst_2  
12: imul  
13: ireturn
```

- **Start:** start bytecode where the variable is visible
- **Length:** number of bytecode bytes during which the variable is visible